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# ON $\alpha$ - AND $\beta$ -RECURSIVELY ENUMERABLE DEGREES

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Several problems in recursion theory on admissible ordinals ( $\alpha$ -recursion theory) and recursion theory on inadmissible ordinals ( $\beta$ -recursion theory) are studied. Fruitful interactions between both theories are stressed. In the first part the admissible collapse is used in order to characterize for some inadmissible  $\beta$  the structure of all  $\beta$ -recursively enumerable degrees as an accumulation of structures of  $\mathfrak{A}$ -recursively enumerable degrees for many admissible structures  $\mathfrak{A}$ . Thus problems about the  $\beta$ -recursively enumerable degrees can be solved by considering "locally" the analogous problem in an admissible  $\mathfrak{A}$  (where results of  $\alpha$ -recursion theory apply). In the second part  $\beta$ -recursion theory is used as a tool in infinite injury priority constructions for some particularly interesting  $\alpha$  (e.g.  $\omega_1^{CK}$ ). New effects can be observed since some structure of the inadmissible world above O' is projected into the  $\alpha$ -recursively enumerable degrees makes it possible to solve some open problems.

A few years ago S.D. Friedman and G.E. Sacks [1] started a new chapter in generalized recursion theory:  $\beta$ -recursion theory. So far recursion theory was studied only on those initial segments  $L_{\alpha}$  of the constructible hierarchy where  $\alpha$  is admissible. In  $\beta$ -recursion theory one considers initial segments  $L_{\beta}$  for any limit ordinal  $\beta$ . This is a natural step since the concept of a recursively enumerable (r.e.) set does not require any closure condition for the considered universe.

The  $\beta$ -r.e. sets are defined to be those subsets of  $L_{\beta}$  which are  $\Sigma_1$ -definable over  $L_{\beta}$ . A function is  $\beta$ -recursive if its graph is  $\beta$ -r.e. Another important concept of recursion theory is finiteness and as in  $\alpha$ -recursion theory a subset of  $L_{\beta}$  is called  $\beta$ -finite if it is an element of  $L_{\beta}$ .

A striking new effect in  $\beta$ -recursion theory is the appearance of  $\beta$ -finite sets which are rather large (compared with the whole universe  $L_{\beta}$ ). If  $\beta$  is inadmissible then there exist  $\beta$ -recursive functions which are cofinal in  $\beta$  and which have as domain a  $\beta$ -finite set. The minimal  $\gamma < \beta$  which occurs as domain of such a function is called the recursive cofinality of  $\beta$  ( $\sigma 1 \text{ cf } \beta$ ). This ordinal is a good measure for the remaining "admissibility" of an inadmissible  $\beta$ . Only  $\beta$ -finite sets of cardinality less than  $\sigma 1 \text{ cf } \beta$  (in  $L_{\beta}$ ) behave like  $\alpha$ -finite sets.

The preceding example shows already that several elementary facts from ordinary recursion theory do not remain true in  $\beta$ -recursion theory. But usually facts of this kind are not considered to be the essential results of recursion theory. Thus the question arises what is happening e.g. with the structure of recursively enumerable degrees. Serious mathematical problems occur here since most of the constructions from  $\alpha$ -recursion theory rely on admissibility and one has to look for new strategies.

In Section 1 of this paper we study  $\beta$ -recursively enumerable degrees for weakly inadmissible  $\beta$  and continue our earlier paper [5]. An inadmissible  $\beta$  is called weakly inadmissible if one can project  $\beta$   $\beta$ -recursively into  $\sigma 1$  cf  $\beta$ , it is called strongly inadmissible otherwise. For weakly inadmissible  $\beta$  one takes a suitable predicate  $\tilde{T} \subseteq L_{\beta}$  which encodes all  $\Delta_0$ -facts about  $L_{\beta}$ . One projects  $\tilde{T}$  by means of the existing  $\beta$ -recursive projection into  $\sigma 1 \operatorname{cf} \beta$ . We write T for the projection of  $\tilde{T}$  and call the structure  $\mathfrak{A} := \langle L_{\sigma 1 \operatorname{cl} \beta}, T \rangle$  the admissible collapse of  $\beta$ . Then every set  $A \subseteq L_{\sigma_1 \text{ cf }\beta}$  is  $\Sigma_1$ -definable over  $L_\beta$  iff it is  $\Sigma_1$ -definable over  $\mathfrak{A}$ . Further  $\mathfrak{A}$  is admissible and therefore we know a lot about  $\mathfrak{A}$ -degrees since most of the numerous results in  $\alpha$ -recursion theory remain true for such an admissible structure with an additional regular predicate (e.g. Shore's density theorem [10] holds for  $\mathfrak{A}$ ). It was shown in [5] that the  $\beta$ -recursive degrees (together with  $\leq_{\alpha}$ ) are isomorphic to the  $\mathfrak{A}$ -r.e. degrees (together with  $\leq_{\mathrm{sr}}$ ). The occurrence of nonzero  $\beta$ -recursive degrees is a typical phenomenon of inadmissible recursion theory. It is easy to see that the greatest  $\beta$ -recursive degree lies strictly between 0 — the degree of the empty set — and the greatest  $\beta$ -r.e. degree 0'. But nothing else is known about nonrecursive r.e. degrees in  $\beta$ -recursion theory. Theorem 1 throws some light onto this problem. A careful analysis shows that we can in fact extend the isomorphism from [5] onto all regular  $\beta$ -r.e. degrees (a set  $A \subseteq L_{\beta}$  is regular if  $\forall \gamma < \beta$  ( $A \cap L_{\gamma} \in L_{\beta}$ ), a  $\beta$ -r.e. degree is regular if it contains a regular  $\beta$ -r.e. set). As isomorphic images we get in  $\mathfrak{A}$  those degrees which are tame r.e. in an  $\mathfrak{A}$ -r.e. degree below it. This leads to the characterization of the structure of all regular  $\beta$ -r.e. degrees as an accumulation of structures of  $\mathfrak{B}$ -r.e. degrees for many different admissible  $\mathfrak{B}$  (see the conclusions following Theorem 1). In particular by applying the splitting theorem to many local structures we get a global splitting theorem for all regular  $\beta$ -r.e. degrees.

Theorem 1 contains in addition a complete characterization of those  $\beta$ -r.e. degrees which are regular. The relative size of  $\sigma 2 \operatorname{cf} \beta$  and  $\sigma 2p\beta$  turns out to be the decisive criterion — a criterion which is well known from  $\alpha$ -recursion theory (see [6] and Shore [12]). In particular for weakly inadmissible  $\beta$  with  $\sigma 2 \operatorname{cf} \beta \ge \sigma 2p\beta$  the regular set theorem from  $\alpha$ -recursion theory holds: every  $\beta$ -r.e. degree is regular. Therefore for these  $\beta$  we get an oversight over all  $\beta$ -r.e. degrees from the isomorphism result.

It should be mentioned that Sack's concept of tameness plays a crucial role in the formulation and the proof of Theorem 1 ( $A \subseteq L_{\beta}$  is defined to be  $\beta$ -tame r.e. if  $\{K \in L_{\beta} \mid K \subseteq A\}$  is  $\beta$ -r.e.).

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Using the result from Theorem 1 we prove in Theorem 2 that the  $\beta$ -recursive degrees are not an initial segment of the  $\beta$ -r.e. degrees for weakly inadmissible  $\beta$ . We do this by performing a suitable construction in the admissible collapse of  $\beta$ .

In Section 2 of the paper we apply results from  $\beta$ -recursion theory in  $\alpha$ -recursion theory. If  $\alpha$  is an admissible ordinal with  $\alpha > \sigma 2$  of  $\alpha \ge \sigma 2p \alpha$ , then the structure of the  $\Sigma_2 L_{\alpha}$  degrees between 0' and 0" is isomorphic to the structure of the  $\mathfrak{B}$ -r.e. degrees for the weakly inadmissible structure  $\mathfrak{B} := \langle L_{\alpha}, C \rangle$  (with a regular  $\alpha$ -r.e. predicate  $C \in 0'$ ). Thus we can apply the preceding results and we find a strange new world of  $\alpha$ -degrees between 0' and 0'' (observe that the case of metarecursion theory is included since  $\omega_1^{CK}$  satisfies the condition  $\alpha > \sigma 2$  cf  $\alpha \ge \sigma 2p \alpha$ ).

The considered  $\alpha$  are of particular interest with respect to the jump because these are the only  $\alpha$  where the jumps of  $\alpha$ -r.e. degrees are not yet known (for  $\Sigma_2$ -admissible  $\alpha$  the situation is exactly as in ordinary recursion theory according to Theorem 5, in the case  $\sigma 2$  cf  $\alpha < \sigma 2p \alpha$  the jump is completely distorted and only the degrees 0',  $0^{\frac{3}{2}}$  and 0" can possibly occur as jumps of  $\alpha$ -r.e. degrees according to [6]). The distinguished degree  $0^{\frac{3}{2}}$  was described in [6] for those  $\alpha$ where incomplete non-hyperregular  $\alpha$ -r.e. degrees exist. Observe that these degrees do always exist in the considered case  $\alpha > \sigma 2$  cf  $\alpha \ge \sigma 2p \alpha$  (an  $\alpha$ -r.e. degree a is hyperregular iff  $\langle L_{\alpha}, A \rangle$  is admissible for regular  $A \in a$ ). It is easy to see that every low degree is hyperregular and every high degree is nonhyperregular (a is low if a' = 0' and a is high if a' = 0''). We show that there are differences among the hyperregular  $\alpha$ -r.e. degree is low and not every nonhyperregular  $\alpha$ -r.e. degree is high.

It is easy to see that  $a' \le 0^{\frac{3}{2}}$  for hyperregular and  $a' \ge 0^{\frac{3}{2}}$  for non-hyperregular  $\alpha$ -r.e. degrees a. Theorem 3 shows that there exist in fact non-hyperregular  $\alpha$ -r.e. degrees a such that  $a' = 0^{\frac{3}{2}}$ . Together with results from [6] we thus get

# $0^{\frac{3}{2}} = \inf \{ a' \mid a \text{ is a non-hyperregular } \alpha \text{-r.e. degree} \}$

for all  $\alpha$  where incomplete non-hyperregular  $\alpha$ -r.e. degrees exists so that we have another characterization of the degree  $0^{\frac{3}{2}}$ . The proof of Theorem 3 is based on a simple trick since the straightforward approach fails. We construct an  $\alpha$ -r.e. set Aas if we want to make it both non-hyperregular and low. This is of course impossible but the constructed non-hyperregular set A is then at least "as low as possible". Thus we get  $A' = 0^{\frac{3}{2}}$ .

Concerning the jump of hyperregular degrees we first observe that there is a rich structure of tame  $\sum_{2} L_{\alpha}$  (or equivalently  $\Delta_{2}L_{\alpha}$ ) degrees between 0' and  $0^{\frac{3}{2}}$  for the considered  $\alpha$  (A is tame  $\sum_{2} L_{\alpha}$  if  $\{K \in L_{\alpha} \mid K \subseteq A\}$  is  $\sum_{2} L_{\alpha}$ ). These  $\alpha$ -degrees are isomorphic to the  $\mathfrak{A}$ -r.e. degrees of an admissible structure  $\mathfrak{A}$ . We show in Theorem 4 that each of these tame  $\sum_{2} L_{\alpha}$  degrees is the jump of a hyperregular  $\alpha$ -r.e. degree. This holds in particular for the greatest tame  $\sum_{2} L$  degree  $0^{\frac{3}{2}}$ . Thus  $0^{\frac{3}{2}}$  is the only point (for any  $\alpha$ ) where the jump of a hyperregular and a non-hyperregular  $\alpha$ -r.e. degree meet together.

We learn from the preceding results that the inadmissible world above 0' casts its shadow upon the central part of  $\alpha$ -recursion theory: the structure of  $\alpha$ -r.e. degrees and sets. In particular we stumble upon the naturally arising notion of an intermediate degree which is characterized by the property  $\mathbf{a}' = 0^{\frac{3}{2}}$ . Essential differences between the structure of r.e. degrees in ordinary recursion theory and the structure of  $\alpha$ -r.e. degrees for some  $\alpha$  with  $\sigma 2 \operatorname{cf} \alpha < \sigma 2p \alpha$  have already been discovered by R. Shore [12]. A further investigation of intermediate degrees may show several differences between ordinary recursion theory and  $\alpha$ -recursion theory for some  $\alpha$  with  $\alpha > \sigma 2 \operatorname{cf} \alpha \ge \sigma 2p \alpha$  (including metarecursion theory). By

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combining recent work of A. Leggett [3] with our results we can give a first example: Martin's celebrated Theorem ("a r.e. degree contains a maximal set iff it is high") is not true in metarecursion theory since for  $\omega_1^{CK}$  there exist maximal sets of intermediate degree.

The combination of Leggett's results with Theorem 4 settles in addition a conjecture from Simpson's thesis [13] positively: There exist in fact hyperregular maximal sets even in nontrivial cases (it turns out that all these sets are of intermediate degree).

In the proof of Theorem 4 we use similar infinite injury strategies as in [6]. But the basic ingredient of the construction is a regular set theorem from  $\beta$ -recursion theory [5] (in fact this is the first application of a regular set theorem from  $\beta$ -recursion theory). All attempts to prove this regular set theorem by using standard methods of  $\alpha$ -recursion theory did fail so far.

Finally in Theorem 5 we show that an admissible  $\alpha$  is  $\Sigma_2$ -admissible iff it satisfies Sack's jump theorem ("every  $\Sigma_2$  degree between 0' and 0" is the jump of an incomplete r.e. degree"). The crucial point of the proof is the demonstration of the failure of Sack's theorem in the case  $\alpha > \sigma 2$  cf  $\alpha \ge \sigma 2p \alpha$ . We apply here another nontrivial result from  $\beta$ -recursion theory (the preceding Theorem 2).

#### **0.** Preliminaries

We use the same notations and definitions as in our preceding papers [5] and [6]. All missing definitions can be found there.

It is not relevant for our arguments whether one chooses  $L_{\beta}$  or  $S_{\beta}$  (see [5]) as the universe for  $\beta$ -recursion theory.

For any structure  $\mathfrak{B} = \langle L_{\beta}, B \rangle$  with  $B \subseteq L_{\beta}$  we say that a subset of  $L_{\beta}$  is  $\mathfrak{B}$ -r.e. ( $\mathfrak{B}$ -recursive) if it is  $\Sigma_1 \mathfrak{B}$  ( $\Delta_1 \mathfrak{B}$ ). We write  $\sigma n$  cf  $\mathfrak{B}$  for the least  $\gamma \leq \beta$  such that a cofinal  $\Sigma_n \mathfrak{B}$  function  $f: \gamma \rightarrow \beta$  exists.  $\sigma np \mathfrak{B}$  is the least  $\gamma \leq \beta$  such that an 1-1  $\Sigma_n \mathfrak{B}$  function  $f: \beta \rightarrow \gamma$  exists. For the special case  $\mathfrak{B} = L_{\beta}$  we write  $\sigma n$  cf  $\beta$  respectively  $\sigma np \beta$ .

Define for any structure  $\mathfrak{B} = \langle L_{\beta}, B \rangle \ \rho_{n,\beta}^{\mathfrak{B}} := \mu \delta \leq \beta$  (a  $\Sigma_n \mathfrak{B}$  set  $M \subseteq \delta$  exists such that  $M \notin L_{\beta}$ ).

The greatest  $\mathfrak{B}$ -recursive degree is always denoted by r.

If  $\mathfrak{B}$  is a weakly inadmissible structure (i.e.  $\beta > \sigma 1$  cf  $\mathfrak{B} \ge \sigma 1p \mathfrak{B}$ ) we reserve the letter  $\kappa$  for  $\sigma 1$  cf  $\mathfrak{B}$  and we write  $\mathfrak{A}$  for the admissible collapse  $\langle L_{\kappa}, T \rangle$  of  $\mathfrak{B}$ . A set  $A \subseteq L_{\kappa}$  is called  $\beta$ -immune if it is immune with respect to neighborhood conditions out of  $L_{\beta} - L_{\kappa}$ , i.e. for any  $K \in L_{\beta}$  we have  $(K \subseteq A \lor K \subseteq L_{\kappa} - A) \rightarrow K \in L_{\kappa}$ . According to [5] one can define for every set  $A \subseteq L_{\kappa}$  a  $\beta$ -immune version  $\tilde{A} \subseteq L_{\kappa}$ of the same  $\mathfrak{A}$ -degree. The operation  $\tilde{P}$  preserves regularity and  $\Sigma_n \mathfrak{A}$ -definability.

For sets  $A, B \subseteq \beta$  we set  $A \lor B := 2A \cup 2B + 1$  where  $2A := \{2x \mid x \in A\}$  and  $2B + 1 := \{2x + 1 \mid x \in B\}$ .

Finally we remind of two conventions. If we write  $L_{\gamma} \models [x \in W]$  for some  $\beta$ -r.e. set W, then this means that we have fixed a  $\Sigma_1 L_{\beta}$  definition  $\phi$  of W and  $L_{\gamma} \models \phi(x)$ .

Further if we write  $L_{\gamma} \models \psi$  for a formula  $\psi$ , then this implies that every parameter in  $\psi$  is an element of  $L_{\gamma}$ .

### **1.** $\beta$ -recursively enumerable degrees

**Theorem 1.** Assume  $\beta$  is weakly inadmissible and **b** is a  $\beta$ -degree. Then (1), (2), (3) are equivalent:

(1) **b** contains a regular  $\beta$ -r.e. set,

(2) **b** contains a  $\beta$ -r.e. set and  $\neg (\mathbf{r} < \mathbf{b} \land \sigma 2 \text{ cf } \beta < \sigma 2p \beta)$  (**r** is the greatest  $\beta$ -recursive degree),

(3) **b** contains a  $\beta$ -immune set  $A \subseteq L_{\kappa}$  such that A is regular over  $L_{\kappa}$  and  $(\mathfrak{A}, B)$ -tame r.e. for some  $\mathfrak{A}$ -r.e. set  $B \subseteq L_{\kappa}$  with  $B \leq_{\mathfrak{A}} A$ .

**Proof.**  $(2) \rightarrow (3)$ . Assume (2) and take a  $\beta$ -r.e. set  $H \in \mathbf{b}$ . We will construct  $\beta$ -immune sets  $B, \hat{A} \subseteq \kappa$  such that  $H =_{\beta} \hat{A}, \langle \mathfrak{A}, B \rangle$  is not strongly inadmissible, B is regular over  $L_{\kappa}, B \leq_{\mathfrak{A}} \hat{A}$  and  $\hat{A}$  is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. This is enough in order to show (3) because we can apply then the regular set theorem for tame r.e. sets in weakly inadmissible structures  $((1) \leftrightarrow (3)$  of Theorem 4 in [5]) respectively the usual regular set theorem for r.e. sets in admissible structures to the  $\langle \mathfrak{A}, B \rangle$ -tame r.e. set  $\hat{A}$  in the weakly inadmissible or admissible structure  $\langle \mathfrak{A}, B \rangle$ . This gives a  $\langle \mathfrak{A}, B \rangle$ -tame r.e. set  $A^*$  with  $A^*$  regular over  $L_{\kappa}$  and  $\hat{A} =_{\langle \mathfrak{A}, B \rangle} A^*$ . Then we take the  $\beta$ -immune set  $\tilde{A}^*$  and define  $A := B \lor \tilde{A}^*$ . It is easy to verify that B, A have all the properties which are demanded in (3).

We construct now the sets  $B, \hat{A}$  with the properties above (i.e. all properties from (3) except  $\hat{A}$  regular).

Fix a  $\beta$ -recursive function P which maps  $L_{\beta}$  1-1 onto  $\kappa$ . Define  $B_1 := P[H]$ .  $B_1$ is  $\mathfrak{A}$ -r.e. and there is a regular  $\mathfrak{A}$ -r.e. set  $B_2$  such that  $B_1 =_{\mathfrak{A}} B_2$ . Then we define  $B := \tilde{B}_2$ . It is obvious that  $B =_{\mathfrak{A}} B_2 =_{\mathfrak{A}} B_1$  and B is  $\mathfrak{A}$ -r.e., regular and  $\beta$ -immune. We further have  $B \leq_{\beta} H$  by the construction of  $B_1$  since B is  $\beta$ -immune and  $B =_{\mathfrak{A}} B_1$ .

Define  $\hat{\mathfrak{A}} := \langle \mathfrak{A}, B \rangle$ . We want to show that  $\hat{\mathfrak{A}}$  is not strongly inadmissible and at this point we use the assumption  $\neg (\mathbf{r} < \mathbf{b} \land \sigma 2 \text{ cf } \beta < \sigma 2p \beta)$  in (2).

(a) *B* is incomplete in  $\mathfrak{A}$ . According to Lemma 3.3 in Shore [10] we have  $\sigma 1 \operatorname{cf}^{\mathfrak{A}} \kappa \ge \rho_{1,\kappa}^{\mathfrak{A}}$ . Further by using the fact that  $\sigma 2p^{\mathfrak{A}} \kappa = \rho_{2,\kappa}^{\mathfrak{A}}$  we get  $\rho_{1,\kappa}^{\mathfrak{A}} \ge \sigma 1p^{\mathfrak{A}} \kappa$  according to Shore [10]. The used equality holds because we have  $\sigma 2p^{\mathfrak{A}} \kappa = \sigma 2p \beta = \rho_{2,\beta} = \rho_{2,\beta}^{\mathfrak{A}}$ . Thus it is proved that  $\sigma 1 \operatorname{cf}^{\mathfrak{A}} \kappa \ge \sigma 1p^{\mathfrak{A}} \kappa$ .

(b) *B* is complete in  $\mathfrak{A}$ . In this case *B* is an element of the  $\beta$ -degree  $\mathbf{r}$  according to Theorem 8 in [5]. We know already that  $B \leq_{\beta} H$  and therefore we have  $\mathbf{r} \leq \mathbf{b}$ . Since (3) follows from Theorem 4 in [5] if  $\mathbf{r} = \mathbf{b}$  we assume  $\mathbf{r} < \mathbf{b}$  for the following. This implies  $\sigma 2 \operatorname{cf} \beta \geq \sigma 2p \beta$  according to the assumption in (2). Therefore we have  $\sigma 1 \operatorname{cf}^{\mathfrak{A}} \kappa = \sigma 2 \operatorname{cf} \beta \geq \sigma 2p \beta = \sigma 1p^{\mathfrak{A}} \kappa$ .

Since  $\hat{\mathfrak{A}}$  is not strongly inadmissible there is a  $\hat{\mathfrak{A}}$ -recursive function  $\hat{P}$  which maps  $\kappa \ 1-1$  onto  $\hat{\kappa} := \sigma 1 \operatorname{cf}^{\hat{\mathfrak{A}}} \kappa$ . Further we fix an  $\hat{\mathfrak{A}}$ -recursive strictly increasing

cofinal function  $\hat{q}: \hat{\kappa} \rightarrow \kappa$ .

We define

$$\begin{split} A_1 &:= \{ \langle 0, x \rangle \in \kappa \mid P^{-1}(x) \cap H \neq \emptyset \} \cup \{ \langle 1, x \rangle \in \kappa \mid P^{-1}(x) \cap L_\beta - H \neq \emptyset \}, \\ A_2 &:= \hat{q}[\hat{P}[A_1]], \qquad A_3 := \tilde{A}_2 \quad \text{and} \quad \hat{A} := B \lor A_3. \end{split}$$

Then  $\hat{A}$  is  $\beta$ -immune by construction and it is obvious that  $B \leq_{\mathfrak{A}} \hat{A}$ . In order to verify the other properties of  $\hat{A}$  we prove:

Claim. Assume that  $K \in L_{\kappa}$ ,  $\kappa$ -card  $(K) < \hat{\kappa}$  and  $K \subseteq A_1 \cap \{1\} \times \kappa$ . Then there exists a set  $K^* \in L_{\kappa}$  such that

$$\kappa$$
-card  $(K^*) < \hat{\kappa}, \qquad P^{-1}[K^*] \subseteq L_{\beta} - H$ 

and

$$\forall \langle 1, x \rangle \in K \quad \exists z \in P^{-1}[K^*](z \in P^{-1}(x))$$

(thus  $K^*$  is a "small"  $\kappa$ -finite set of witnesses for " $K \subseteq A_1 \cap \{1\} \times \kappa$ ").

Proof of the claim. We have  $P[H] = B_1 =_{\mathfrak{A}} B$ . Therefore we can write

$$y \in \kappa - P[H] \wedge P^{-1}(y) \in P^{-1}(x)$$

as a  $\Sigma_1 \hat{\mathfrak{A}}$  formula  $\psi(x, y)$ . Since  $K \subseteq A_1 \cap \{1\} \times \kappa$  we have that

 $\forall x \in K \quad \exists y \in \kappa(\hat{\mathfrak{A}} \models \psi(x, y)).$ 

Since  $\kappa$ -card  $(K) < \hat{\kappa}$  there exists in fact a  $\kappa$ -finite function  $h: K \to \kappa$  such that  $\forall x \in K(\hat{\mathfrak{A}} \models \psi(x, h(x)))$  (show first that the function  $g: K \to \kappa$  such that

 $\forall x \in K(g(x) = \mu \delta(\langle L_{\delta}, T \cap L_{\delta}, B \cap L_{\delta}) \models [\exists y \psi(x, y)])$ 

is  $\kappa$ -finite because  $\kappa$ -card  $(K) < \hat{\kappa} = \sigma 1 \operatorname{cf}^{\hat{\mathfrak{A}}} \kappa$ ).

Then the set  $K^* := \operatorname{Rg} h$  has all the properties which are demanded in the claim.  $\Box$ 

We can show now that  $\hat{A}$  is  $\hat{\mathfrak{A}}$ -tame r.e.: We have

$$K \in L_{\kappa} \land K \subseteq \hat{A} \leftrightarrow \exists \hat{K} \in L_{\kappa} (\langle K, \hat{K} \rangle \in W_{e} \land \kappa \text{-card} (\hat{K}) < \hat{\kappa} \land \hat{K} \subseteq A_{1} \cap \{1\} \times \kappa \}$$

for some suitable  $\hat{\mathfrak{A}}$ -r.e. set  $W_e$ . By the claim we further have

$$\begin{split} \hat{K} &\in L_{\kappa} \wedge \kappa \operatorname{-card} (\hat{K}) < \hat{\kappa} \wedge \hat{K} \subseteq A_{1} \cap \{1\} \times \kappa \\ \leftrightarrow \exists K^{*} \in L_{\kappa} (\kappa \operatorname{-card} (K^{*}) < \hat{\kappa} \wedge P^{-1}[K^{*}] \subseteq L_{\beta} - H \\ \wedge \forall \langle 1, x \rangle \in \hat{K} \quad \exists z \in P^{-1}[K^{*}] (z \in P^{-1}(x))). \end{split}$$

The latter can be written as a  $\Sigma_1 \hat{\mathfrak{A}}$  formula since

 $P^{-1}[K^*] \subseteq L_{\beta} - H \leftrightarrow K^* \subseteq L_{\kappa} - B_1$ 

and  $B_1 \leq_{\mathfrak{N}} B$ . Thus we have shown that  $\hat{A}$  is  $\hat{\mathfrak{A}}$ -tame r.e.

 $H \leq_{\beta} \hat{A}$  follows from

$$K \in L_{\beta} \land K \subseteq H \leftrightarrow \exists x \in \kappa (P^{-1}(x) = K \land \langle 1, x \rangle \notin A_1),$$
  
$$K \in L_{\beta} \land K \subseteq L_{\beta} - H \leftrightarrow \exists x \in \kappa (P^{-1}(x) = K \land \langle 0, x \rangle \notin A_1)$$

and  $A_1 \leq_{w\mathfrak{A}} \hat{A}$  (it is obvious that  $A_1 \leq_{w\mathfrak{A}} \hat{A}$ ; this implies  $A_1 \leq_{w\mathfrak{A}} \hat{A}$  because  $B \leq_{\mathfrak{A}} \hat{A}$ ).

It only remains to show that  $\hat{A} \leq_{\beta} H$ . We have

$$K \in L_{\kappa} \land K \subseteq A_{2} \Leftrightarrow K \in L_{\kappa} \land \exists \hat{K} \in L_{\kappa} (K = \hat{q}[\hat{P}[\hat{K}]] \land \kappa \text{-card} (\hat{K}) < \hat{\kappa})$$
$$\land \exists K_{0}K_{1}(K = K_{0} \cup K_{1} \land K_{0} \subseteq \{0\} \times \kappa \cap A_{1} \land K_{1} \subseteq \{1\} \times \kappa \cap A_{1}).$$

The part  $K_0 \subseteq \{0\} \times \kappa \cap A_1$  is obviously  $\Sigma_1 L_\beta$  since  $\beta$ -card  $(K_0) < \kappa$ . Further we have shown before that for  $\kappa$ -finite sets  $K_1$  of  $\kappa$ -cardinality less than  $\hat{\kappa}$  we can write  $K_1 \subseteq \{1\} \times \kappa \cap A_1$  as a  $\Sigma_1 \hat{\mathfrak{A}}$  formula. Since  $B \leq_\beta H$  we have altogether expressed  $K \in L_\kappa \wedge K \subseteq A_2$   $\beta$ -recursively in H. Since  $\hat{A}$  is  $\beta$ -immune we can therefore express  $K \in L_\beta \wedge K \subseteq \hat{A}$   $\beta$ -recursively in H.

In order to show the other part of  $\hat{A} \leq_{\beta} H$  we observe that

$$K \in L_{\kappa} \land K \subseteq \kappa - A_1 \leftrightarrow \exists K_1 K_2 \in L_{\beta}(K_1 = \bigcup \{\check{K} \mid \langle 0, P(\check{K}) \rangle \in K\}$$
$$\land K_1 \subseteq L_{\beta} - H \land K_2 = \bigcup \{\check{K} \mid \langle 1, P(\check{K}) \rangle \in K\} \land K_2 \subseteq H).$$

We further have

 $K \in L_{\kappa} \wedge K \subseteq L_{\kappa} - A_2 \leftrightarrow \exists \hat{K} \in L_{\kappa}(\hat{q}[\hat{P}[\hat{K}]] = \operatorname{Rg} \hat{q} \cap K \wedge \hat{K} \subseteq \kappa - A_1).$ 

The right side of this equivalence can be written as a  $\Sigma_1 \hat{\mathfrak{A}}$  formula since  $\hat{q}$  is strictly increasing and continuous. If we combine these facts it is easy to see that  $K \in L_{\beta} \wedge K \subseteq L_{\beta} - \hat{A}$  can be expressed  $\beta$ -recursively in H.

 $(3) \rightarrow (1)$ . Take sets A, B according to (3). One can assume that B is in addition  $\beta$ -immune and regular over  $L_{\kappa}$ . We are going to define a  $\Pi_1 L_{\beta}$  set H which is regular over  $L_{\beta}$  such that  $A =_{\beta} H$ . Then  $L_{\beta} - H$  will be a regular  $\beta$ -r.e. set of the same  $\beta$ -degree as A.

One might try to understand the definition of H as follows: There exists some  $\Pi_1 \mathfrak{A}$  set D such that  $w \in A \Leftrightarrow \exists z \in L_{\kappa}(\langle w, z \rangle \in D)$ . H is some sort of  $\Pi_1 \mathfrak{A}$ -uniformization of this relation D.

First we define a set  $\hat{H}$ . Fix a  $\Sigma_1(\mathfrak{A}, B)$  formula  $\phi$  such that  $R \subseteq A \leftrightarrow \langle \mathfrak{A}, B \rangle \models \phi(R)$ . For  $w \in A$  we put a 4-tupel  $\langle w, \gamma, \delta, K \rangle$  into  $\hat{H}$  such that

(a)  $\gamma$  is minimal such that

$$\langle L_{\gamma}, T \cap L_{\gamma}, B \cap L_{\gamma} \rangle \models \exists R(w \in R \land \phi(R)),$$

(b)  $K = B \cap L_{\gamma}$ ,

(c)  $\delta \ge \gamma$  is minimal such that

 $\forall x \in K(\langle L_{\delta}, T \cap L_{\delta} \rangle \models \psi(x))$ 

where  $\psi$  is some fixed  $\Sigma_1 \mathfrak{A}$  definition of *B*.

For every  $w \in A$  there exists exactly one tripel  $\langle \gamma, \delta, K \rangle$  such that  $\langle w, \gamma, \delta, K \rangle \in \hat{H}$  and  $\hat{H}$  is  $\Pi_1 \mathfrak{A}$  definable.

Fix a  $\beta$ -finite function U which maps  $L_{\kappa}$  1–1 onto  $\kappa$  and a  $\beta$ -recursive strictly increasing and continuous cofinal function q from  $\kappa$  into  $\beta$ . We write  $\tilde{q}$  for  $q \circ U$ .

Define the set H as follows (we use set theoretic pairing):

$$\langle x, \gamma, \delta, K \rangle \in H : \leftrightarrow x \in \tilde{q}[L_{\kappa}] \land \langle \tilde{q}^{-1}(x), \gamma, \delta, K \rangle \in H.$$

H is  $\Pi_1 L_{\beta}$  definable since  $\tilde{q}[L_{\kappa}]$  is  $\Delta_1 L_{\beta}$  and  $\hat{H}$  is  $\Pi_1 L_{\beta}$ .

Take some set  $R_0 \in L_{\kappa}$  such that  $R_0 \subseteq A$ . Then there exists some  $\gamma_0 < \kappa$  such that for all  $w \in R_0$ 

$$\langle L_{\gamma_0}, T \cap L_{\gamma_0}, B \cap L_{\gamma_0} \rangle \models \exists R(w \in R \land \phi(R))$$

(simply choose  $\gamma_0$  such that  $R_0 \in L_{\gamma_0}$  and  $\langle L_{\gamma_0}, T \cap L_{\gamma_0}, B \cap L_{\gamma_0} \rangle \models \phi(R_0)$ ). It follows from this observation that  $\hat{H} \upharpoonright R_0$  is  $\beta$ -finite. We further have that  $\tilde{q} \upharpoonright R_0$  is  $\beta$ -finite and therefore  $H \upharpoonright \tilde{q}[R_0]$  is  $\beta$ -finite as well. Observe that the given argument relies mainly on the fact that A is  $\langle \mathfrak{A}, B \rangle$ -tame r.e.

We can now prove that H is regular over  $L_{\beta}$ . Assume that some  $\gamma$  such that  $\kappa \leq \gamma < \beta$  is given. In order to show that  $L_{\gamma} \cap H \in L_{\beta}$  we define

$$K := \{ w \in L_{\kappa} \mid \exists y_1 y_2 y_3 \in L_{\kappa}(\langle \tilde{q}(w), y_1, y_2, y_3 \rangle \in L_{\gamma}) \}.$$

By using the properties of  $\tilde{q}$  we see that  $K \in L_{\kappa}$ . Since A is regular over  $L_{\kappa}$  we have  $R := K \cap A \in L_{\kappa}$ . By the preceding we have then  $H \upharpoonright \tilde{q}[R] \in L_{\beta}$ . Since  $H \cap L_{\gamma} = H \upharpoonright \tilde{q}[R] \cap L_{\gamma}$  this implies that  $H \cap L_{\gamma} \in L_{\beta}$ .

 $A \leq_{\scriptscriptstyle B} H$ . We have

 $R \in L_{\beta} \land R \subseteq A \leftrightarrow R \in L_{\kappa} \land \exists F \in L_{\beta}((F \text{ is a function}) \land \text{dom } F = \tilde{q}[R] \land F \subseteq H)$ and

$$R \in L_{\beta} \land R \subseteq L_{\kappa} - A \leftrightarrow R \in L_{\kappa} \land R \subseteq L_{\kappa} - A \leftrightarrow \tilde{q}[R] \times L_{\kappa}^{3} \subseteq L_{\beta} - H$$

where  $\tilde{q}[R] \times L^3_{\kappa}$  is a  $\beta$ -finite set due to the properties of set theoretic pairing.

 $H \leq_{\beta} A$ . We write  $\pi_0, \pi_1, \pi_2, \pi_3$  for the projection functions which are associated with the set theoretic 4-tuples. We have:

$$\begin{split} F &\in L_{\beta} \wedge F \subseteq H \Leftrightarrow F \in L_{\beta} \wedge (F \text{ is a function}) \wedge \exists R \in L_{\kappa} (\text{dom } F \\ &= \tilde{q}[R] \wedge R \subseteq A \wedge \exists \gamma \delta K \in L_{\kappa} (\gamma \leqslant \delta < \kappa \wedge \pi_{1}[F] \subseteq \gamma \wedge K \subseteq L_{\gamma} \\ &\wedge L_{\gamma} - K \subseteq L_{\gamma} - B \wedge \exists K_{1}(K_{1} = L_{\delta} \cap T \wedge \forall x \in K(\langle L_{\delta}, K_{1} \rangle \models \psi(x)) \\ &\wedge \pi_{2}[F] \subseteq \delta \wedge \forall \hat{K} \in \pi_{3}[F] \exists \hat{\gamma} < \gamma(\hat{K} = K \cap L_{\hat{\gamma}}) \wedge \exists \rho < \beta(\kappa < \rho \\ &\wedge F \in L_{\rho} \wedge \tilde{q} \upharpoonright R \in L_{\rho} \wedge L_{\rho} \models [F \text{ is correctly defined with respect to} \\ \tilde{q} \upharpoonright R, \gamma, \delta, K, K_{1} \text{ according to the definition of } H])))). \end{split}$$

Since T is  $\Delta_1 L_{\beta}$  and  $B \leq_{\beta} A$  we can express in this way  $F \in L_{\beta} \land F \subseteq H \beta$ -recursively in A.

Concerning the other part of  $H \leq_{\beta} A$  we have

$$K \in L_{\beta} \land K \subseteq L_{\beta} - H \leftrightarrow \exists K_1 K_2 K_3 F \in L_{\beta} (K_1 = \pi_0 [K] \cap \tilde{q} [L_{\kappa}]$$
$$\land K_2 \cup K_3 = \tilde{q}^{-1} [K_1] \land K_2 \subseteq A \land K_3 \subseteq L_{\kappa} - A$$
$$\land \operatorname{dom} F = \tilde{q} [K_2] \land F \subseteq H \land F \cap K = \emptyset).$$

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The expression " $F \subseteq H$ " which occurs on the right side can be expressed  $\beta$ -recursively in A as we have shown before. Observe that the equivalence relies on the fact that H is a function.

(1) $\rightarrow$ (3). Assume  $H \subseteq L_{\beta}$  is regular and  $\beta$ -r.e. We construct sets A, B with the properties in (3).

Fix a  $\beta$ -recursive function P which maps  $L_{\beta}$  1-1 onto  $\kappa$ . Define B := P[H]. Then B is  $\mathfrak{A}$ -r.e. and regular over  $L_{\kappa}$ .

Define

$$A_1 := \{ \langle 0, x \rangle \in \kappa \mid P^{-1}(x) \cap H \neq \emptyset \} \cup \{ \langle 1, x \rangle \in \kappa \mid P^{-1}(x) \cap L_\beta - H \neq \emptyset \}$$

and  $A := \tilde{A}_1$ . It is obvious that  $H \leq_{\beta} A$  and A is  $\beta$ -immune. Further the regularity of H over  $L_{\beta}$  implies the regularity of  $A_1$  and A over  $L_{\kappa}$ .

A is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. We have for any  $K \in L_{\kappa}$ :

$$\begin{split} K &\subseteq A_1 \leftrightarrow \exists K_1 K_2 \in L_{\kappa} (K = K_1 \cup K_2 \wedge K_1 \subseteq \{0\} \times \kappa \wedge K_2 \subseteq \{1\} \times \kappa \\ & \wedge \exists \gamma < \beta(L_{\gamma} \models [\forall \langle 0, x \rangle \in K_1 \exists y \in P^{-1}(x)(y \in H)]) \\ & \wedge \exists \hat{K} \in L_{\kappa} (\hat{K} \subseteq \kappa - B \wedge \forall \langle 1, x \rangle \in K_2 \exists y \in P^{-1}[\hat{K}](y \in P^{-1}(x)))). \end{split}$$

Whereas " $\leftarrow$ " is obvious we have to show for " $\rightarrow$ " that the set  $\hat{K} \in L_{\kappa}$  exists. The set  $\bigcup \{P^{-1}(x) \mid \langle 1, x \rangle \in K_2\}$  is  $\beta$ -finite and therefore subset of some  $L_{\gamma}$  with  $\gamma < \beta$ . Since  $L_{\gamma} \cap H \in L_{\beta}$  there exists a  $\beta$ -finite function h such that

$$\forall \langle 1, x \rangle \in K_1(h(\langle 1, x \rangle) \in P^{-1}(x) \cap L_{\gamma} - H).$$

Then the  $\kappa$ -finite set  $\hat{K} := P[\operatorname{Rg} h]$  has all the desired properties. We have thus shown that  $A_1$  is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. which implies that A is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. as well.

The equivalence above proves simultaneously one part of  $A \leq_{\beta} H$  if we write  $P^{-1}[\hat{K}] \subseteq L_{\beta} - H$  instead of  $\hat{K} \subseteq \kappa - B$  on the right side. The other part is immediate since A is  $\beta$ -immune (see the analogous reduction in the proof of  $(2) \rightarrow (3)$ ).

Since  $H \leq_{\beta} A$  and B = P[H] it is obvious that  $B \leq_{\mathfrak{A}} A$ .

 $(3) \rightarrow (2)$ . Assume for a contradiction that  $r < b \land \sigma 2$  cf  $\beta < \sigma 2p \beta$ . Then we can choose the sets A, B according to (3) such that in addition  $B <_{\mathfrak{A}} A$  and  $B \in 0'$  in  $\mathfrak{A}$  is regular. In this case the structure  $\mathfrak{A} := \langle \mathfrak{A}, B \rangle$  is strongly inadmissible because

$$\sigma 1 \operatorname{cf}^{\mathfrak{A}} \kappa = \sigma 2 \operatorname{cf}^{\mathfrak{A}} \kappa = \sigma 2 \operatorname{cf} \beta < \sigma 2p \beta = \sigma 2p^{\mathfrak{A}} \kappa = \sigma 1p^{\mathfrak{A}} \kappa.$$

We have shown in Lemma 24 in [5] that  $A \leq_{\beta} \emptyset$  for every regular tame r.e. set  $A \subseteq L_{\beta}$  if  $\beta$  is strongly inadmissible. The argument works as well for our strongly inadmissible *structure*  $\hat{\mathfrak{A}}$  since by Jensen's Uniformization Theorem [2] we have  $\sigma 2p \beta = \rho_{2,\beta}$  so that  $\sigma 1 \operatorname{cf}^{\hat{\mathfrak{A}}} \kappa < \rho_{2,\beta} = \rho_{\hat{\mathfrak{A}}}^{\hat{\mathfrak{A}}}$ . This inequality  $\sigma 1 \operatorname{cf}^{\hat{\mathfrak{A}}} \kappa < \rho_{\hat{\mathfrak{A}},\kappa}^{\hat{\mathfrak{A}}}$  is needed for the argument. Therefore we have  $A \leq_{\hat{\mathfrak{A}}} \emptyset$  which implies that  $A \leq_{\hat{\mathfrak{A}}} B$ . This is a contradiction to the assumption  $B <_{\hat{\mathfrak{A}}} A$  so that we have proved that  $\neg (\mathbf{r} < \mathbf{b} \land \sigma 2 \operatorname{cf} \beta < \sigma 2p \beta)$  if  $\mathbf{b}$  contains a set according to (3). It remains to show that  $\mathbf{b}$  contains a  $\beta$ -r.e. set but this follows from (3)  $\rightarrow$  (1).

Thus the proof of Theorem 1 is complete.

## Conclusions

(1) Consider a weakly inadmissible  $\beta$  with admissible collapse  $\mathfrak{A}$ . According to Theorem 8 in [5] there exists an isomorphism I from  $\langle S, \leq_{\mathfrak{A}} \rangle$  onto  $\langle R, \leq_{\beta} \rangle$  where S is the set of  $\mathfrak{A}$ -r.e.  $\mathfrak{A}$ -degrees and R is the set of  $\beta$ -recursive (or equivalently  $\beta$ -tame r.e.)  $\beta$ -degrees.

Define  $\hat{S}$  as the set of all  $\mathfrak{A}$ -degrees which contain a regular set A such that A is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. for some  $\mathfrak{A}$ -r.e. set B with  $B \leq_{\mathfrak{A}} A$ . Define  $\hat{R}$  as the set of all regular  $\beta$ -r.e.  $\beta$ -degrees. Then we have: There exists an isomorphism  $\hat{I}$  from  $\langle \hat{S}, \leq_{\mathfrak{A}} \rangle$  onto  $\langle \hat{R}, \leq_{\beta} \rangle$  such that  $\hat{I} \upharpoonright S = I$ .

The definition of  $\hat{I}$  is simple: If  $a \in \hat{S}$  is an  $\mathfrak{A}$ -degree then we take a set  $A \in a$  with the properties as in the definition of  $\hat{S}$  such that A is in addition  $\beta$ -immune (a contains such an A by the properties of the  $\tilde{}$ -operation in [5]). Define then  $\hat{I}(a)$  as the  $\beta$ -degree of A. Theorem 1 implies that the so defined function  $\hat{I}$  is an isomorphism.

(2) For the first time we now have an overview over the structure of all  $\beta$ -r.e.  $\beta$ -degrees for an inadmissible  $\beta$ . Assume  $\beta$  is weakly inadmissible and  $\sigma 2$  cf  $\beta \ge \sigma 2p \beta$ .

In this case the set  $\hat{R}$  in the preceding conclusion is the set of all  $\beta$ -r.e.  $\beta$ -degrees. Further all structures  $\langle \mathfrak{A}, B \rangle$  which occur in the definition of  $\hat{S}$  for this case are not strongly inadmissible. Therefore we can drop the requirement "A regular" in the definition of  $\hat{S}$  (apply the usual regular set theorem for admissible structures respectively the regular set theorem for tame r.e. sets in weakly inadmissible structures which is contained in Theorem 4 in [5]).

Thus we learn that there are many  $\beta$ -r.e. degrees between  $\mathbf{r}$  and 0' in  $\beta$ : The structure of all  $\beta$ -r.e. degrees  $\mathbf{b}$  such that  $\mathbf{r} \leq \mathbf{b} \leq 0'$  (together with  $\leq_{\beta}$ ) is isomorphic to the structure of all  $\hat{\mathfrak{A}}$ -r.e.  $\hat{\mathfrak{A}}$ -degrees (together with  $\leq_{\hat{\mathfrak{A}}}$ ) for an admissible structure  $\hat{\mathfrak{A}}$ .

We get  $\hat{\mathfrak{A}}$  by applying the admissible collapse two times if necessary: It is easy to see that  $\hat{I}$  maps exactly the tame  $\Sigma_2 \mathfrak{A}$   $\mathfrak{A}$ -degrees  $a \ge 0_{\mathfrak{A}}'$  onto the considered  $\beta$ -degrees b, because in the definition of an  $\mathfrak{A}$ -degrees  $a \in \hat{S}$  we can always assume that  $B \in 0'$  if a satisfies  $0' \le a$  ( $\mathfrak{A}$  is as always the admissible collapse of  $\beta$ ). Fix then an  $\mathfrak{A}$ -r.e. regular set  $C \in 0_{\mathfrak{A}}'$ . The structure  $\langle \mathfrak{A}, C \rangle$  is either admissible, in which case we define  $\hat{\mathfrak{A}} := \langle \mathfrak{A}, C \rangle$ , or it is weakly inadmissible, in which case we define  $\hat{\mathfrak{A}}$  to be the admissible collapse of  $\langle \mathfrak{A}, C \rangle$ .

In order to describe all  $\beta$ -r.e. degrees we write RE  $(\hat{\mathfrak{A}})$  for the structure of all  $\hat{\mathfrak{A}}$ -r.e. degrees together with  $\leq_{\hat{\mathfrak{A}}}$  for any admissible structure  $\hat{\mathfrak{A}}$ . For the considered  $\beta$  we can describe then the structure  $\langle \hat{R}, \leq_{\beta} \rangle$  of all  $\beta$ -r.e. degrees as an accumulation of many structures RE  $(\mathfrak{A}_b)$  where **b** ranges over the  $\beta$ -tame r.e. degrees. The picture we get is familiar from fireworks: Every  $\beta$ -tame r.e. degree **b** is the starting point of some structure RE  $(\mathfrak{A}_b)$  above **b**, where **b** itself corresponds to  $0_{\mathfrak{A}_c}$ .

To be a little more exact, we start with the structure  $\langle \hat{R}, \leq_{\beta} \rangle$  and go then to the isomorphic structure  $\langle \hat{S}, \leq_{\mathfrak{A}} \rangle$ . Consider a degree  $\boldsymbol{a} \in \hat{S}$ . Then there is a set  $A \in \boldsymbol{a}$ 

which is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. for some  $\mathfrak{A}$ -r.e. B with  $B \leq_{\mathfrak{A}} A$ . It is obvious that we can choose B in addition  $\beta$ -immune so that the  $\beta$ -degree of B is some  $\beta$ -tame r.e. degree **b**. In order to define  $\mathfrak{A}_b$  we first observe that the structure of the  $\langle \mathfrak{A}, B \rangle$ -tame r.e.  $\mathfrak{A}$ -degrees  $a \leq_{\mathfrak{A}} B$  (together with  $\leq_{\mathfrak{A}}$ ) is isomorphic to the structure of the  $\langle \mathfrak{A}, B \rangle$ -tame r.e.  $\langle \mathfrak{A}, B \rangle$ -degrees (together with  $\leq_{\mathfrak{A}, B}$ ). We define  $\mathfrak{A}_b$  as the structure  $\langle \mathfrak{A}, B \rangle$ , if  $\langle \mathfrak{A}, B \rangle$  is admissible. Otherwise  $\langle \mathfrak{A}, B \rangle$  is weakly inadmissible and we define  $\mathfrak{A}_b$  to be the admissible collapse of  $\langle \mathfrak{A}, B \rangle$ .

For those weakly inadmissible  $\beta$  which satisfy  $\sigma 2 \operatorname{cf} \beta < \sigma 2p \beta$  one can give a similar description except that we don't know which structure has to be attached to the greatest  $\beta$ -tame r.e. degree r as an characterization of the  $\beta$ -r.e.  $\beta$ -degrees between r and 0'. We expect that the degree structure of a strongly inadmissible structure occurs at this point.

Thus we see that for all  $\beta$ -r.e.  $\beta$ -degrees d > 0 such that

$$\neg (\mathbf{r} < \mathbf{d} \land \sigma 2 \operatorname{cf} \beta < \sigma 2p \beta)$$

the splitting theorem holds: There exist  $\beta$ -r.e.  $\beta$ -degrees  $d_1$ ,  $d_2$  such that  $0 < d_1 < d_1$ ,  $0 < d_2 < d$  and d is the least upper bound of  $d_1$  and  $d_2$ .

The claim is immediate from [5] if d is  $\beta$ -tame r.e. Otherwise we apply the splitting theorem for r.e. degrees in admissible structures to the admissible structure  $\mathfrak{A}_b$  in which d is represented by a r.e. degree.

We do not yet know much about the "overlapping" of the structures RE  $(\mathfrak{A}_b)$ . Concerning a proof of the density theorem for regular  $\beta$ -r.e. degrees one can eliminate this problem<sup>1</sup>. The following trick in ORT is due to David Posner. Consider sets  $A_1, A_2 \subseteq \omega$  such that  $A_2 <_{\omega} A_1$  and there exist r.e. sets  $B_1, B_2$  such that  $A_i$  is r.e. in  $B_i$  and  $B_i \leq_{\omega} A_i$ , i = 1, 2. In order to find sets  $A_3, B_3$  such that  $A_2 <_{\omega} A_3 <_{\omega} A_1, B_3 \leq_{\omega} A_3$  and  $A_3$  is r.e. in  $B_3$  one considers the following cases:

(a)  $B_1 \lor A_2 = {}_{\omega}A_2$ . Apply the density theorem for r.e. degrees in  $\langle L_{\omega}, B_1 \lor B_2 \rangle$ .

(b)  $B_1 \lor A_2 = {}_{\omega}A_1$ . Then  $A_2 < {}_{\omega}B_1 \lor A_2$  and  $B_1 \lor A_2$  is r.e. in  $B_2$ . Apply the density theorem for r.e. degrees in  $\langle L_{\omega}, B_2 \rangle$ .

(c)  $A_2 \leq_{\omega} B_1 \lor A_2 \leq_{\omega} A_1$ . Define  $A_3 := B_1 \lor A_2$ .

By combining this argument with Theorem 1 and the density theorem in  $\alpha$ -RT (Shore [10]) one sees that the regular  $\beta$ -r.e. degrees are dense if  $\beta$  is weakly inadmissible and  $\sigma 1$  cf  $\beta$  is a cardinal in L.

For a general weakly inadmissible  $\beta$  the problem is reduced to the density theorem for r.e. degrees in admissible structures  $\langle L_{\alpha}, B \rangle$  with B regular. The latter problem is open, even if B is in addition  $\alpha$ -r.e. (the proof of this case might suffice for our application). Shore's proof of the density theorem for  $\alpha$ -r.e. degrees [10] uses properties of projecta which are dubious in presence of a predicate B.

(3) One has to be careful in generalizing the preceding results to weakly inadmissible structures  $\mathfrak{B} = \langle L_{\beta}, B \rangle$ . We have used that  $\sigma 2p \beta = \rho_{2,\beta}$  (uniformization theorem [2]) at two points in the proof of Theorem 1 (in (2) $\rightarrow$ (3) and in

<sup>&</sup>lt;sup>1</sup>We are grateful to Sy D. Friedman and Gerald E. Sacks who informed us about this. Further we would like to thank Richard A. Shore and the referee for pointing out the situation concerning relativized projecta.

 $(3) \rightarrow (2)$ ). Even for some admissible  $\mathfrak{B}$  this equality is false. Fortunately the equality holds for the most interesting applications in  $\alpha$ -recursion theory:

If  $\alpha$  is admissible,  $\alpha > \sigma 2$  cf  $\alpha \ge \sigma 2p \alpha$  and  $B \in 0'$  is  $\alpha$ -r.e. and regular over  $L_{\alpha}$  then we have for  $\mathfrak{B} := \langle L_{\beta}, B \rangle$  that

 $\sigma 2p^{\mathfrak{B}} \alpha = \sigma 3p \alpha = \rho_{3,\alpha} = \rho_{2,\alpha}^{\mathfrak{B}}$ 

by the uniformization theorem for  $\alpha$ . (Observe that the level 3 is the first one where we really need the full power of the uniformization theorem in  $\alpha$ -recursion theory since  $\Sigma_2$ -uniformization is trivial for admissible  $\alpha$ .) Thus we get a lot of information about  $\Sigma_2 L_{\alpha}$  degrees above 0' for these  $\alpha$ .

**Theorem 2.** Assume that  $\beta$  is weakly inadmissible. Then the  $\beta$ -tame r.e. degrees are not an initial segment of the  $\beta$ -r.e. degrees.

**Proof.** According to Theorem 1 (and Theorem 4 in [5]) it is enough to solve the following problem for the admissible collapse  $\mathfrak{A} = \langle L_{\kappa}, T \rangle$  of  $L_{\beta}$ : Construct a set A such that  $A <_{\mathfrak{A}} 0_{\mathfrak{A}}'$ , A is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. for some  $\mathfrak{A}$ -r.e. set B with  $B \leq_{\mathfrak{A}} A$ and there exists no  $\mathfrak{A}$ -r.e. set W such that  $A =_{\mathfrak{A}} W$ .

We will solve this problem in a way which was suggested by M. Lerman for the special case of ordinary recursion theory. It turns out that routine precautions are sufficient to make the proof work for all admissible structures.

Call a set D  $\mathfrak{A}$ -d.r.e. if D = A - B for some  $\mathfrak{A}$ -r.e. sets A, B. It is easy to see that we have in this case  $D \leq_{\mathfrak{A}} 0'_{\mathfrak{A}}$  if A or B is regular over  $L_{\kappa}$ .

**Lemma 1.** Every  $\mathfrak{A}$ -d.r.e. set D is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. for some  $\mathfrak{A}$ -r.e. set B with  $B \leq_{\mathfrak{A}} A$ .

**Proof.** Assume that  $D = A_1 - B_1$  with  $\mathfrak{A}$ -r.e. sets  $A_1, B_1$ . We can assume without loss of generality that  $B_1 \subseteq A_1$ . Fix an  $\mathfrak{A}$ -recursive 1-1 enumeration  $f: \kappa \to A_1$  of the  $\mathfrak{A}$ -r.e. set  $A_1$ . Define  $B := f^{-1}[B_1]$ . It is obvious that B is  $\mathfrak{A}$ -r.e. We have for every  $K \in L_{\kappa}$ :

$$K \subseteq D \leftrightarrow K \subseteq A_1 - B_1 \leftrightarrow \exists \hat{K} \in L_{\kappa}(f[\hat{K}] = K \land \hat{K} \subseteq \kappa - B).$$

This shows that D is  $\langle \mathfrak{A}, B \rangle$ -tame r.e. Finally we have  $B \leq_{\mathfrak{A}} D$  because  $K \subseteq \kappa - B \Leftrightarrow f[K] \subseteq D$ .

In order to prove Theorem 2 it is thus enough to construct an  $\mathfrak{A}$ -d.r.e. set D = A - B with A regular such that D is not contained in any  $\mathfrak{A}$ -r.e. degree. We will mainly describe the construction since the verification of the desired properties is fairly standard for this finite injury priority construction. For convenience in writing we restrict our attention to an admissible set  $L_{\alpha}$  instead of an admissible structure  $\mathfrak{A}$  (for which the proof is litterally the same).

 $D_{\sigma}, A_{\sigma}, B_{\sigma}, W_{a,\sigma}$  are the collections of elements which have been put into these sets  $D, A, B, W_a$  before step  $\sigma \in \alpha$  of the construction.

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For every  $e = \langle a, b, c \rangle \in \alpha$  with  $a, b, c \in \alpha$  we have a requirement

$$R_e := (\neg W_a \leq^b_\alpha D \lor \neg D \leq^c_{w\alpha} W_a).$$

(we have fixed a universal  $\Sigma_1 L_{\alpha}$  predicate  $U_1$  and we write  $W_{\alpha}$  for  $\{x \mid \langle x, a \rangle \in U_1\}$ , we have further fixed an enumeration of  $U_1$  so that the notation  $W_{a,\sigma}$  makes sense).

Define  $\rho := t\sigma 2p \alpha$  (the tame  $\Sigma_2$  projectum of  $\alpha$ , see [4]) and fix a tame  $\Sigma_2$  projection g which maps  $\rho$  1-1 onto  $\alpha$ . Further we fix an  $\alpha$ -recursive tame approximation  $f: \alpha \times \rho \to \alpha$  such that  $f(\sigma, \cdot)$  is 1-1 for every  $\sigma \in \alpha$ . We further define an approximation  $\{c\}_{\sigma}^{W_{\alpha,\sigma}}$  for every function  $\{c\}_{w_{\alpha}}^{W_{\alpha}}$  (this is the — in general partial — function which is weakly  $\alpha$ -recursive in  $W_{\alpha}$  with index c). Define

$$\begin{aligned} \{c\}_{\sigma}^{W_{a,\sigma}}(x) \downarrow : \Leftrightarrow \exists \tau \leq \sigma \exists KHy \in L_{\tau}(L_{\tau} \models [\langle x, y, K, H \rangle \in W_{c}] \\ \land K \subseteq W_{a,\tau} \land H \subseteq L_{\sigma} - W_{a,\sigma}). \end{aligned}$$

Observe that the  $\sigma$  in  $W_{a,\sigma}$  at the end is not a misprint. If  $\{c\}_{\sigma}^{W_{a,\sigma}}(x)\downarrow$  we determine the value such that  $\{c\}_{\sigma}^{W_{a,\sigma}}(x)\simeq y$  and the negative neighborhood of this computation as follows: Choose  $\tau \leq \sigma$  with the properties above minimal. For this  $\tau$  choose  $\hat{y}$ ,  $\hat{K}$ ,  $\hat{H}$  with the properties above such that  $\langle \hat{y}, \hat{K}, \hat{H} \rangle$  is minimal with respect to the canonical  $\Delta_1 L_{\alpha}$  well-ordering of  $L_{\alpha}$ . Then we define  $\{c\}_{\sigma}^{W_{a,\sigma}}(x)\simeq \hat{y}$  and  $\hat{H}$  is defined to be the negative neighborhood of this computation.

Every requirement is at every step  $\sigma$  of the construction in one of the states 0, 1, 2. At the beginning of the construction every requirement is in state 0.

We say that  $R_e$  requires attention at step  $\sigma$  if  $\exists \delta < \rho(f(\sigma, \delta) = e)$  and

(1)  $R_e$  is in state 0 at the beginning of step  $\sigma$  and there exists some  $x < \sigma$  such that

(a)  $x - A_{\sigma}$  has an order type  $\geq \gamma$  where  $\gamma := \sup f(\sigma, \cdot)[\delta + 1]$ , and

(b)  $x \notin A_{\sigma}$ , and

(c) if at some step  $\sigma' < \sigma$  a requirement  $R_{e'}$  received attention with  $f(\sigma', \delta') = e'$  and  $\delta' < \delta$  then we have  $\sigma' \leq x$  and

(d)  $\{c\}_{\sigma^{a,o}}^{W_{a,o}}(x) \approx 0$  with a negative neighborhood K (we assume that the characteristic function of a set has value 1 for arguments *in* the set), and

(e) we have for K from (d) that

$$\exists H_1 H_2 \in L_{\sigma}(L_{\sigma} \models [\langle K, 1, H_1, H_2 \rangle \in W_b] \land H_1 \subseteq D_{\sigma} \land H_2 \subseteq L_{\sigma} - D_{\sigma})$$

(see the definition of  $W_a \leq^b_{\alpha} D$ ); or

(2)  $R_e$  was put into state 1 at some step  $\sigma' < \sigma$  and an element x was put into A at step  $\sigma'$  and the state of  $R_e$  was not changed after  $\sigma'$  and we have  $\{c\}_{\sigma}^{W_{\alpha,\sigma}}(x) \simeq 1$ ; or

(3)  $R_e$  is not in state 0 at the beginning of step  $\sigma$  and there exists no  $\sigma' < \sigma$  and  $\delta < \rho$  such that  $f(\tau, \delta) \simeq e$  for all  $\tau$  such that  $\sigma' \leq \tau \leq \sigma$ .

### **Construction**

Step  $\sigma$ . If no requirement requires attention at step  $\sigma$  go to the next step. Otherwise choose  $\delta < \rho$  minimal such that  $R_e$  with  $e = f(\sigma, \delta)$  requires attention at step  $\sigma$ . We say then that  $R_e$  receives attention at step  $\sigma$  and we do the following:

At first we put all  $R_{\tilde{e}}$  into state 0 such that  $\tilde{e} = f(\sigma, \tilde{\delta})$  for some  $\tilde{\delta} > \delta$ . Then we proceed as follows

Case 1.  $R_e$  requires attention according to (1).

Choose the x in (1) minimal and put into A. Further we put  $R_e$  into state 1.

Case 2.  $R_e$  requires attention according to (2) but not according to (3).

Put the x in (2) into B and put  $R_e$  into state 2.

Case 3.  $R_e$  requires attention according to (3).

Put  $R_e$  into state 0.

End of the construction.

One proves as usual that for every  $\delta < \rho$  there exists a step such that no  $R_e$  with  $g^{-1}(e) \leq \delta$  receives attention after this step.

The condition " $x - A_{\sigma}$  has an order type  $\ge \sup f(\sigma, \cdot)[\delta + 1]$ " in (1) (a) makes it possible to show that  $\alpha - A$  has order type  $\alpha$  and A is regular over  $L_{\alpha}$ .

Assume then for a contradiction that  $D =_{\alpha} W_a$  for some  $\alpha$ -r.e. set  $W_a$ . It is essential that we can assume without loss of generality that  $W_a$  is regular (apply the regular set theorem). Assume that b, c are indices such that  $W_a \leq_{\alpha}^{b} D$  and  $D \leq_{w\alpha}^{c} W_a$  (it is important that we take b such that  $W_a \leq_{\alpha}^{b} D$ —not just  $W_a \leq_{w\alpha}^{b} D$ ). We consider then requirement  $R_e$  with  $e := \langle a, b, c \rangle$ .

Consider a step  $\sigma_0$  such that after step  $\sigma_0$  no requirement of higher priority than  $R_e$  requires attention and such that  $x \in \sigma_0 - A$  exists so that x - A has an order type which is large enough to satisfy the condition in (1) (a) for  $R_e$  at  $\sigma_0$ . A step with these properties exists by our previous remarks.

Since  $W_a$  is regular,  $x \notin D$  and  $D \leq_{w\alpha}^c W_a$  we have that  $\{c\}_{\sigma}^{W_{\alpha,\sigma}}(x) \approx 0$  with the same negative neighborhood K for all large enough  $\sigma$  and we have  $K \subseteq L_{\alpha} - W_{\alpha}$ .

Since  $W_a \leq_{\alpha}^{b} D$  there exists a step  $\sigma_1 \geq \sigma_0$  and there exist neighborhoods  $H_1, H_2 \in L_{\sigma_1}$  such that  $H_1 \subseteq D_{\sigma}$  for all  $\sigma \geq \sigma_1$  and  $H_2 \subseteq L_{\alpha} - D_{\sigma}$  for  $\sigma \geq \sigma_1$  (we have used for the latter that A is regular).

It follows from these observations that there is some step  $\hat{\sigma}$  such that  $R_e$  is put into state 1 at step  $\hat{\sigma}$ , some  $\hat{x}$  is put into A at step  $\hat{\sigma}$  (call the associated negative neighborhood  $\hat{K}$ , thus  $\hat{K} \subseteq L_{\hat{\sigma}} - W_{a,\hat{\sigma}}$ ) and  $R_e$  is not put into state 0 at any step after  $\hat{\sigma}$ .

Then either  $R_e$  is in state 1 at all steps after  $\hat{\sigma}$  — in which case we have  $\hat{x} \in A - B =: D$  but not  $\{c\}^{W_a}(x) \simeq 1$ , a contradiction.

Or there exists a first step  $\check{\sigma} > \hat{\sigma}$  such that  $R_e$  is put into state 2 at step  $\check{\sigma}$ ,  $\hat{x}$  is put into B at step  $\check{\sigma}$  and  $R_e$  is in state 2 at all steps after  $\check{\sigma}$ . In this case we have that some element  $y \in \hat{K}$  was enumerated into  $W_a$  at some step between  $\hat{\sigma}$  and  $\check{\sigma}$ although some computation $\langle \hat{K}, 1, \hat{H}_1, \hat{H}_2 \rangle \in W_b$  with  $\hat{H}_1 \subseteq D_{\hat{\sigma}} \land H_2 \subseteq L_{\hat{\sigma}} - D_{\hat{\sigma}}$  existed at step  $\hat{\sigma}$ . This computation may have been injured (if  $\hat{x} \in \hat{H}_2$ ) at the end of step  $\hat{\sigma}$  by putting  $\hat{x}$  into A, but this  $\hat{x}$  is the only possible injury of the computation. Since one puts  $\hat{x}$  into  $B \subseteq L_{\alpha} - D$  at step  $\check{\sigma}$  the computation is restored in any case at step  $\check{\sigma}$ . Further this computation in D will remain valid at all steps after  $\check{\sigma}$  so that we have in fact  $\hat{H}_1 \subseteq D \land \hat{H}_2 \subseteq L_{\alpha} - D$ . This contradicts our assumption  $W_a \leq_{\alpha}^b D$  since we had observed before that  $\neg \hat{K} \subseteq L_{\alpha} - W_a$ .

This finishes the proof of Theorem 2.

### 2. α-recursively enumerable degrees

Not much is proved so far about the jump of  $\alpha$ -r.e. degrees if  $\alpha$  is  $\Sigma_2$  admissible or  $\alpha > \sigma 2$  cf  $\alpha \ge \sigma 2p \alpha$  (these are the types (1) and (2) in the classification of Section 3 in [6]), whereas everything is known about the jump of  $\alpha$ -r.e. degrees for the other  $\alpha$  (see [6]). We will study in this chapter mainly type (2). The result for type (1) is contained in Theorem 5.

For  $\alpha$  of type (2) we know so far that there exist nonzero  $\alpha$ -r.e. degrees which are low and incomplete  $\alpha$ -r.e. degrees which are high. This agrees with the situation in ordinary recursion theory. On the other hand we learnt in the meantime that the structure of  $\Sigma_2 L_{\alpha}$  degrees above 0' is for this type very different from the corresponding structure in ordinary recursion theory: There exists the distinguished degree  $0^{\frac{3}{2}}$  between 0' and 0" which was described in Lemma 7, Section 2 in [6]. Observe that for  $\alpha$  of type (2) there exist always incomplete non-hyperregular  $\alpha$ -r.e. degrees so that the assumption of this Lemma is satisfied. We can apply the results about weakly inadmissible structures from [5] and the preceding Section 1. The structure  $\mathfrak{B} := \langle L_{\alpha}, C \rangle$  where  $C \in 0'$  is  $\alpha$ -r.e. and regular is weakly inadmissible if  $\alpha$  is of type (2) and it is obvious that the  $\mathfrak{B}$ -r.e. degrees are isomorphic to the  $\Sigma_2 L_{\alpha}$  degrees above 0' in  $L_{\alpha}$ .

It is easy to see that  $0^{\frac{3}{2}}$  is the boundary between the jumps of hyperregular  $\alpha$ -r.e. degrees and the jumps of non-hyperregular  $\alpha$ -r.e. degrees for  $\alpha$  of type (2):

If **a** is a hyperregular  $\alpha$ -r.e. degree, then **a**' is a tame  $\Sigma_2 L_{\alpha}$  degree and therefore we have  $\mathbf{a}' \leq_{\alpha} 0^{\frac{3}{2}}$  because  $0^{\frac{3}{2}}$  is the greatest tame  $\Sigma_2 L_{\alpha}$  degree.

If **a** is a non-hyperregular  $\alpha$ -r.e. degree then we have that the complete  $\Sigma_2 L_{\alpha}$  set  $U_2$  is weakly  $\alpha$ -recursive in **a'** (Shore [11]). According to Lemma 7 in [6] we have therefore  $0^{\frac{3}{2}} \leq_{\alpha} a'$ .

In particular we have thus shown that  $0^{\frac{3}{2}}$  is comparable with the jump of every  $\alpha$ -r.e. degree. But we do not yet know so far whether there exist for  $\alpha$  of type (2) any other jumps of  $\alpha$ -r.e. degrees besides 0' and 0".

**Theorem 3.** Assume  $\alpha$  is admissible and  $\alpha > \sigma 2$  cf  $\alpha \ge \sigma 2p \alpha$ . Then there exists a non-hyperregular  $\alpha$ -r.e. degree **a** such that  $\mathbf{a}' = 0^{\frac{3}{2}}$ .

**Proof.** One runs into a lot of trouble if one tries to prove this Theorem as one would do it in ordinary recursion theory, i.e. if one fixes a set  $S \in 0^{\frac{3}{2}}$  and tries to construct A as a suitable "thick subset" as in Soare [14] in order to get  $A' = \alpha 0^{\frac{3}{2}}$ .

Therefore we prove Theorem 3 as follows: We make sure that the order type of  $\alpha - A$  is less than  $\alpha$  so that A is non-hyperregular as in Shore [9]. This implies that  $0^{\frac{3}{2}} \leq_{\alpha} A'$ .

We keep the jump down in order to get  $0^{\frac{3}{2}} = {}_{\alpha}A'$  by following the usual strategy which is applied to make a constructed set low (see e.g. Soare [14] Theorem 4.1.). This strategy will make sure that " $K \in L_{\alpha} \wedge K \subseteq A'$ " is  $\Pi_2 L_{\alpha}$  which implies  $A' \leq_{\alpha} 0^{\frac{3}{2}}$  by the special properties of  $0^{\frac{3}{2}}$ .

For the exact proof we fix a tame  $\Sigma_2 L_{\alpha}$  projection P from  $L_{\alpha}$  onto  $t\sigma 2p \alpha = \sigma 2$  cf  $\alpha =: \kappa$  and an  $\alpha$ -recursive tame approximation  $P(\cdot, \cdot): \alpha \times L_{\alpha} \rightarrow \alpha$  such that  $P(\sigma, \cdot)$  is 1-1 for every  $\sigma \in \alpha$ .

Further we fix a cofinal strictly increasing and continuous  $\Sigma_2 L_{\alpha}$  function  $g: \kappa \to \alpha$  and an  $\alpha$ -recursive approximation  $g(\cdot, \cdot): \alpha \times \kappa \to \alpha$  such that  $g(\sigma, \cdot) \leq \sigma$  and  $g(\sigma, \cdot)$  is weakly increasing. According to the definition of the jump in  $\alpha$ -recursion theory (see Shore [11]) there exists an  $\alpha$ -r.e. set W such that for every  $M \subseteq L_{\alpha}$ :

$$M' = \{ y \mid \exists H_1 H_2(\langle y, H_1, H_2 \rangle \in W \land H_1 \subseteq M \land H_2 \subseteq L_\alpha - M) \}.$$

We fix an  $\alpha$ -recursive enumeration of W such that  $W_{\sigma} \subseteq L_{\sigma}$  for every  $\sigma$ .

We define a restriction function  $r(i, \sigma)$  for arguments  $i \in \kappa$  and  $\sigma \in \alpha$ .  $r(i, \sigma)$  will be the  $\alpha$ -finite set of those elements less than  $\sigma$  which are kept out from A at step  $\sigma$  with priority *i*.

Fix *i* and  $\sigma$  for the following definition of  $r(i, \sigma)$ . Let *K* be the set of those elements  $x \in \sigma - A_{\sigma}$  such that  $x - A_{\sigma}$  has an order type less than *i*. Further we check for every j < i whether the following condition (\*j) is satisfied:

$$\exists \tau \leq \sigma \exists y (\forall \tau' (\tau \leq \tau' \leq \sigma \rightarrow P(\tau', y) = j) \land \exists H_1 H_2 (\langle y, H_1, H_2 \rangle)$$
  
$$\in W_\tau \land H_1 \subseteq A_\tau \land L_\tau \models [\operatorname{card} (H_2) < \kappa] \land H_2 \subseteq L_\tau - A_\sigma)).$$

If (\*j) is satisfied we choose  $\tau$  in this condition minimal. For this  $\tau \text{ let } \langle y, \hat{H}_1, \hat{H}_2 \rangle$  be the minimal tripel (with respect to  $\langle L_a \rangle$ ) which satisfies (\*j). We define then  $K_i := \hat{H}_2$ . If (\*j) is not satisfied we define  $K_j := \emptyset$ . Then we define

 $r(i, \sigma) := K \cup \bigcup \{K_i \mid j < i\}.$ 

Since  $\kappa$  is a regular  $\alpha$ -cardinal we have  $\alpha$ -card  $(r(i, \sigma)) < \kappa$  for every  $i, \sigma$ .

We have positive requirements  $P_i$  for every  $i < \kappa$  which try to make sure that the order type of g(i) - A is less than  $\kappa$ .

Construction. Step  $\sigma$ . Choose  $i < \kappa$  minimal such that

$$g(\sigma, i) - A_{\sigma} - r(i, \sigma) \neq \emptyset.$$

If such an *i* does not exist go to the next step. Otherwise for this *i* we say that  $P_i$  receives attention at  $\sigma$ . We put all elements of  $g(\sigma, i) - A_{\sigma} - r(i, \sigma)$  into A.

End of the construction.

Claim 1. For all  $i < \kappa$  there exists a step  $\sigma_i$  such that no  $P_j$  with  $j \le i$  receives attention after step  $\sigma_i$ .

**Proof.** Assume  $i_0$  is minimal such that  $\sigma_{i_0}$  does not exist. Since  $i_0 < \sigma^2$  cf  $\alpha$  there exists a step  $\sigma'$  such that no  $P_i$  with  $j < i_0$  receives attention after  $\sigma'$ .

There exists  $\hat{\sigma} \ge \sigma'$  such that

$$\forall \sigma \geq \hat{\sigma} (\forall i \leq i_0 (g(\sigma, i) = g(i) < \sigma) \land \forall y \in L_{\alpha} ((P(y) \leq i_0)))$$
  
$$\rightarrow P(\sigma, y) = P(y)) \land (P(y) > i_0 \rightarrow P(\sigma, y) > P(y))).$$

Choose  $\tau_1 > \hat{\sigma}$  minimal such that  $P_{i_0}$  receives attention at step  $\tau_1$ . Then all elements of  $g(i_0) - A_{\tau_1} - r(i_0, \tau_1)$  are put into A at step  $\tau_1$ .

Choose  $\tau_2 > \tau_1$  minimal such that  $P_{i_0}$  receives attention at  $\tau_2$ . This implies that  $g(i_0) - A_{\tau_2} - r(i_0, \tau_2) \neq \emptyset$ . That is only possible if there exists some  $y \in r(i_0, \tau_1)$  such that  $y \notin r(i_0, \tau_2) \land y \notin A_{\tau_2}$ . Then there exists some  $j < i_0$  such that  $y \in K_i$  where  $K_i$  is some negative neighborhood  $\hat{H}_2 \subseteq L_{\tau_1} - A_{\tau_1}$  (see the definition of  $r(i_0, \tau_1)$ ) and such that a minimal step  $\tau$  with  $\tau_1 < \tau < \tau_2$  exists at which some  $z \in \hat{H}_2$  is put into A. By the choice of  $\tau_1, \tau_2$  some requirement  $P_i$  with  $i > i_0$  must then receive attention at  $\tau$ . But since  $\hat{H}_2 \subseteq \tau - A_{\tau}$  one has for this  $i, \tau$  that  $\hat{H}_2 \subseteq r(i, \tau)$  so that no element of  $\hat{H}_2$  is put into A at step  $\tau$ . Contradiction.

Claim 2.  $\alpha - A$  is unbounded in  $\alpha$  and has order type  $\kappa$ . Further A is regular and non-hyperregular.

**Proof.** (a) Assume for a contradiction that some  $\gamma$  less than  $\kappa$  is the order type of  $\alpha - A$ .

Go to a stage  $\sigma$  such that after  $\sigma$  no  $P_i$  with  $i \leq \gamma$  receives attention. Consider the minimal *j* such that requirement  $P_j$  receives attention at some step  $\tau > \sigma$ . We have then  $g(\tau, j) - A_{\tau} - r(j, \tau) \neq \emptyset$  and therefore  $\tau - A_{\tau}$  has an order type  $\geq j$  by the definition of  $r(j, \tau)$ . The first *j* elements of  $\tau - A_{\tau}$  will never be put into *A* by the choice of  $\tau > \sigma$ . This is a contradiction to  $j > \gamma$ .

(b) Assume for a contradiction that some  $\delta < \alpha$  exists such that  $\delta - A$  has order type  $\kappa$ .

Choose *i* such that  $g(i) > \delta$  and  $\sigma$  such that no  $P_j$  with  $j \le i$  receives attention at  $\sigma$  and  $g(\sigma, i) = g(i)$ .

Then we have  $g(\sigma, i) - A_{\sigma} - r(i, \sigma) \neq \emptyset$  because  $r(i, \sigma)$  has  $\alpha$ -cardinality less than  $\kappa$  and  $g(\sigma, i) - A$  has order type  $\kappa$ . Therefore some  $P_j$  with  $j \leq i$  receives attention at  $\sigma$  which is a contradiction to the choice of  $\sigma$ .

A is regular because  $\delta - A$  has an order type less than  $\kappa = 2 \operatorname{cf} \alpha$  for every  $\delta < \alpha$ . A is non-hyperregular because the function  $f: \kappa \to \alpha$ , where f(i) is the *i*-th element of  $\alpha - A$ , is cofinal and weakly  $\alpha$ -recursive in A.

Claim 3.  $A' =_{\alpha} 0^{\frac{3}{2}}$ .

**Proof.** We have

$$K \subseteq A' \leftrightarrow \forall i < \kappa \forall \sigma \forall y ((y \in K \land P(y) = i \land \forall \tau \ge \sigma ((\text{no } P_i \\ \text{with } j \le i \text{ receives attention at } \tau) \land P(\tau, y) = i))$$
  
$$\rightarrow \exists \tau \ge \sigma \exists H_1 H_2(\langle y, H_1, H_2 \rangle \in W_\tau \land H_1 \subseteq A_\tau \land H_2 \\ \subseteq \tau - A_\tau \land L_\tau \models [\text{card } (H_2) < \kappa])).$$

In order to prove this equivalence we use for " $\rightarrow$ " that A is regular and that  $\delta - A$  has an order type less than  $\kappa$  for every  $\delta < \alpha$ . For " $\leftarrow$ " we consider some  $y \in K$ . Then there exist i = P(y) and  $\sigma$  such that the premise of the right side is satisfied. Therefore there exists  $\tau \ge \sigma$  as in the conclusion on the right side. Thus there exists some tripel  $\langle y, \hat{H}_1, \hat{H}_2 \rangle \in W$  such that  $\hat{H}_1 \subseteq A_{\tau}$ ,  $\hat{H}_2 \subseteq \tau - A_{\tau}$  and  $\hat{H}_2 \subseteq r(j, \tau')$  for every  $\tau' \ge \tau$  and j > i. Therefore no element of  $\hat{H}_2$  will be put into A at any step  $\tau' \ge \tau$  because no  $P_j$  with  $j \le i$  receives attention after  $\tau \ge \sigma$ . This shows that  $y \in A'$ .

The right side of the equivalence above can obviously be written in  $\Pi_2 L_{\alpha}$  form so that we have

$$K \in L_{\alpha} \land K \subseteq A' \leftrightarrow \langle K, e \rangle \notin U_2$$

for some fixed index *e* where  $U_2$  is a universal  $\Sigma_2 L_{\alpha}$  predicate. Since we have  $U_2 \leq_{w\alpha} 0^{\frac{3}{2}}$  by Lemma 7 in [6] we have thus expressed  $K \subseteq A' \alpha$ -recursively in  $0^{\frac{3}{2}}$ .

Since  $K \in L_{\alpha} \wedge K \subseteq L_{\alpha} - A'$  is as well  $\prod_{2} L_{\alpha}$  (trivial) we have shown that  $A' \leq_{\alpha} 0^{\frac{3}{2}}$ .

Finally  $0^{\frac{3}{2}} \leq_{\alpha} A'$  follows from the non-hyperregularity of A.

This finishes the proof of Theorem 3.

**Theorem 4.** Assume  $\alpha$  is admissible and  $\alpha > \sigma 2$  cf  $\alpha \ge \sigma 2p \alpha$ . If D is a tame  $\Sigma_2 L_{\alpha}$  set with  $0' \le_{\alpha} D$ , then there exists a hyperregular  $\alpha$ -r.e. set A such that  $A' =_{\alpha} D$ .

**Proof.** The general strategy will be that one which is used in ordinary recursion theory in order to prove Sack's jump theorem [7] (see Soare [14]). This strategy is a variation of the strategy which is used in order to construct incomplete high r.e. sets. The positive requirements are the same but instead of using Sack's preservation strategy one tries here to keep the jump down by preserving computations which predict that some element is going to be in the jump of the constructed set. This is the same preservation strategy as in the construction of non-zero low degrees and we have used this strategy already in the proof of the preceding Theorem.

If one wants to transfer this construction from ordinary recursion theory to  $\alpha$ -recursion theory one has to overcome similar problems as in the construction of incomplete high  $\alpha$ -r.e. degrees (Theorem 1 in [6]). But most of these problems have to be solved here in a different way because of slight differences in the situation.

We have again the difficulty that the construction from ordinary recursion theory gives only  $D \leq_{w\alpha} A'$  instead of  $D \leq_{\alpha} A'$ . Here we have to keep the constructed set A hyperregular and therefore we use the regular set theorem from  $\beta$ -recursion theory. According to Theorem 4 in [5] there exists a tame  $\Sigma_2 L_{\alpha}$  set  $\hat{D}$ in the  $\alpha$ -degree of D which is regular and satisfies in addition for every set  $B \subseteq L_{\alpha}: D \leq_{w\alpha} B \leftrightarrow D \leq_{\alpha} B$ . Observe that this escape was not possible in the construction of incomplete high  $\alpha$ -r.e. degrees because 0" does not contain a set with the latter property if  $\alpha$  is of type (2).

In Theorem 1 in [6] we could get along without a regular representative in 0" because we made the priority list extremely short. For the present proof one needs a better approximation to the priority list than there because here we have the additional requirement to keep the jump down. Thus it is a lucky circumstance that due to the regular set theorem from  $\beta$ -recursion theory we can work here with a regular representative and use an — in general longer — priority list of length  $\sigma^2$  cf  $\alpha = t\sigma^2 p \alpha$ .

At limit points of the priority list we have here the same problem with the inductive argument as in [6]. This problem was described there in point 4) of the motivation before Theorem 1. Similar as there we use a fixpoint argument in order to get along although the induction hypothesis is too weak at limit points.

For the exact proof consider the weakly inadmissible structure  $\mathfrak{B} := \langle L_{\alpha}, C \rangle$  with  $C \in 0' \alpha$ -r.e. and regular. Since D is  $\mathfrak{B}$ -tame r.e. there exists by Theorem 1 in [5] a regular  $\mathfrak{B}$ -tame r.e. set  $S \subseteq \alpha$  such that

$$\forall B \subseteq L_{\alpha}(S \leq_{w\mathfrak{B}} B \leftrightarrow S \leq_{\mathfrak{B}} B), \qquad D = \mathfrak{B}S$$

and every  $K \subseteq S$  with  $K \in L_{\alpha}$  has an order type less than  $\sigma 2 \operatorname{cf} \alpha$  (we need the latter property for the proof of Claim 1).

We define  $\hat{D} := C \lor S$  and have then  $D =_{\alpha} \hat{D}$ ,  $\hat{D}$  is regular and tame  $\Sigma_2 L_{\alpha}$  and  $\hat{D} \leq_{w\alpha} B \leftrightarrow \hat{D} \leq_{\alpha} B$  for all  $B \subseteq L_{\alpha}$  such that  $0' \leq_{\alpha} B$ .

Fix  $\Delta_0$ -formulas  $\phi$ ,  $\hat{\psi}$  such that

$$x \in C \leftrightarrow L_{\alpha} \models \exists y \phi(x, y)$$

and

$$x \in S \leftrightarrow L_{\alpha} \models \exists y \forall z ((y \text{ is an ordinal}) \land \psi(x, y, z)).$$

Define

$$\psi(x, y, z) := (y \text{ is an ordinal}) \land \hat{\psi}(x, y, z).$$

We will construct A as a "thick subset" of the  $\alpha$ -r.e. set  $R \subseteq \alpha$  which is defined as follows:

$$\langle u, v \rangle \in R \leftrightarrow u, v \in \alpha \land (\exists x (v = 2x \land \neg L_u \models \exists y \phi(x, y)))$$
$$\lor \exists x (v = 2x + 1 \land \forall y \leq u \exists z \neg \psi(x, y, z))).$$

We have

$$v \notin \hat{D} \leftrightarrow \{u \mid \langle u, v \rangle \in R\} = \alpha,$$
  
$$v = 2x \land x \in C \leftrightarrow \{u \mid \langle u, v \rangle \in R\} = \mu \gamma (L_{\gamma} \models \exists y \phi(x, y)) \in \alpha$$

and

$$v = 2x + 1 \land x \in S \leftrightarrow \{u \mid \langle u, v \rangle \in R\} = \mu \gamma (\forall z \psi(x, \gamma, z)) \in \alpha.$$

We write in the following for any set  $M \subseteq \alpha$ :

$$M^{(v_0)} := \{ \langle u, v \rangle \in M \mid v = v_0 \}$$

and

$$M^{K} := \bigcup \{ M^{(v)} \mid v \in K \} \text{ for any set } K.$$

Claim 1. For every set  $\hat{K} \in L_{\alpha}$  there exist  $K, H \in L_{\alpha}$  such that  $R^{\hat{K}} = H \cup \alpha \times K$ . In particular  $R^{\hat{K}}$  is  $\alpha$ -recursive.

**Proof.** Define  $K_1, K_2 \in L_{\alpha}$  by  $\hat{D} \cap \hat{K} = K_1 \vee K_2$  (we use here the regularity of  $\hat{D}$ ).

Define a function  $f: K_1 \to \alpha$  such that  $f(x) = \mu \gamma(L_\gamma \models \exists y \phi(x, y))$ . Then f is  $\alpha$ -finite because f is  $\alpha$ -recursive and dom  $f = K_1 \in L_\alpha$ . Therefore  $R^{2K_1}$  is some  $\alpha$ -finite set  $H_1$  (we had defined  $2K_1 := \{2x \mid x \in K_1\}$ ).

Further by the choice of  $\hat{D}$  respectively S we have that  $\alpha$ -card  $(K_2) < \sigma 2$  cf  $\alpha$ . Since  $R^{(v)}$  is  $\alpha$ -finite for every  $v \in K_2$  we get that  $R^{2K_2+1}$  is some  $\alpha$ -finite set  $H_2$ .

We have then  $R^{\hat{K}} := H \cup \alpha \times K$  with the  $\alpha$ -finite sets  $H := H_1 \cup H_2$ ,  $K := \hat{K} - \hat{D} \cap \hat{K}$ .  $\Box$ 

Fix as in the proof of Theorem 3 an  $\alpha$ -r.e. set W such that

 $M' = \{ \mathbf{y} \mid \exists H_1 H_2 \in L_{\alpha}(\langle \mathbf{y}, H_1, H_2 \rangle \in W \land H_1 \subseteq M \land H_2 \subseteq \alpha - M \}.$ 

We assume here for trivial technical reasons that  $M' \subseteq \alpha$  (use some  $\alpha$ -recursive function which maps  $L_{\alpha}$  1-1 onto  $\alpha$ ).

We fix  $\alpha$ -recursive enumerations of the sets W and R such that  $W_{\sigma} \subseteq L_{\sigma}$  and  $R_{\sigma} \subseteq L_{\sigma}$  for all  $\sigma \in \alpha$ .

As usual we write  $A_{\sigma}$ ,  $W_{\sigma}$  etc. for the set of those elements which are enumerated *before* step  $\sigma$ . Further we write  $R_{\sigma}^{(v)}$ ,  $R_{\sigma}^{K}$  instead of  $R^{(v)} \cap R_{\sigma}$ respectively  $R^{K} \cap R_{\sigma}$ .

The restriction function r will be defined in two parts. First we define a restriction function q which is needed in order to keep the jump A' down. Then we define a restriction function  $\hat{r}$  which is needed in order to make A hyperregular (we use the standard strategy in order to make A hyperregular).

For  $\gamma, \sigma \in \alpha$  we define  $q(\gamma, \sigma)$  as follows: Check whether some step  $\tau < \sigma$  exists such that

$$\exists H_1 H_2(\langle \gamma, H_1, H_2 \rangle \in W_{\tau} \land H_1 \subseteq A_{\tau} \land H_2 \subseteq \tau - A_{\sigma}).$$

If  $\tau$  does not exist define  $q(\gamma, \sigma) := 0$ . Otherwise we take the least such  $\tau$  and we take for this  $\tau$  the least  $z \in \alpha$  such that the existing sets  $H_1, H_2$  can chosen to be subsets of z. We define then

$$q(\gamma, \sigma) := \max\{z, 1\}.$$

In order to define  $\hat{r}$  we first have to define analogously as in Soare [14] the two

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functions  $l: \alpha \times \alpha \rightarrow \alpha$  and  $u: \alpha \times \alpha \times \alpha \rightarrow \alpha$ . Define

$$l(e, \sigma) := \begin{cases} \min(\{x \mid \{(e)_0\}_{\sigma'}^{A_{\sigma'}}(x) \uparrow\} \cup \{\alpha^*, (e)_1\}) & \text{if } \alpha^* < \alpha, \\ \min(\{x \mid \{(e)_0\}_{\sigma'}^{A_{\sigma'}}(x) \uparrow\} \cup \{(e)_1\}) & \text{if } \alpha^* = \alpha. \end{cases}$$

We use in this definition the two projection functions  $(\cdot)_0$  and  $(\cdot)_1$  which are associated with the pairing function  $\langle \cdot, \cdot \rangle : \alpha \times \alpha \to \alpha$ . The approximations  $\{e\}^{A}_{\sigma}(x)$ to the functions which are weakly recursive in A are defined as follows: We write  $\{e\}^{A}_{\sigma}(x) \downarrow$  if

$$\exists \tau < \sigma \exists y H_1 H_2(L_{\tau} \models [\langle x, y, H_1, H_2 \rangle \in W_e] \land H_1 \subseteq A_{\tau} \land H_2 \subseteq L_{\tau} - A_{\sigma}).$$

The "use function" u is defined as follows: We set  $u(e, x, \sigma) = 0$  if  $\{e\}_{\sigma}^{A_{\sigma}}(x) \uparrow$ . If  $\{e\}_{\sigma}^{A_{\sigma}}(x) \downarrow$  we go back to the definition of  $\{e\}_{\sigma}^{A_{\sigma}}(x) \downarrow$  and choose the existing  $\tau < \sigma$  minimal. For this  $\tau$  we take the minimal tripel  $\langle \hat{y}, \hat{H}_1, \hat{H}_2 \rangle$  which satisfies the condition in the definition. We define then  $u(e, x, \sigma)$  as the minimal  $z \in \alpha$  such that  $\hat{H}_2 \subseteq z$ . The  $\hat{y}$  out of this minimal tripel is defined to be the value of  $\{e\}_{\sigma}^{A_{\sigma}}(x)$  and we write then  $\{e\}_{\sigma}^{A_{\sigma}}(x) \simeq \hat{y}$ .

Finally we define

$$\hat{r}(e, \sigma) := \sup \left\{ u((e)_0, x, \sigma) \mid x < l(e, \sigma) \right\}$$

and

$$r(e, \sigma) := \max(\hat{r}(e, \sigma), q(e, \sigma)).$$

Observe that we always have  $r(e, \sigma) \leq \sigma$ .

In order to assign priorities we fix a tame  $\Sigma_2 L_{\alpha}$  projection P which maps  $\alpha$  1-1 onto  $\kappa := \sigma 2$  cf  $\alpha$ . We fix an  $\alpha$ -recursive tame approximation  $P(\cdot, \cdot) : \alpha \times \alpha \to \alpha$ such that  $P(\sigma, \cdot)$  is 1-1 for every  $\sigma \in \alpha$  and

$$\forall \gamma < \kappa \exists \sigma \forall \sigma' \ge \sigma \forall z ((P(z) \le \gamma \rightarrow P(\sigma', z) = P(z)) \\ \wedge (P(z) > \gamma \rightarrow P(\sigma', z) > \gamma)).$$

Construction. Step  $\sigma$ . Consider every  $x = \langle x', e \rangle \in R_{\sigma+1}$  which is not already element of  $A_{\sigma}$ . We put x into A at step  $\sigma$  if and only if  $x \ge r(i, \sigma)$  for every  $i \le \sigma$  such that  $P(\sigma, i) \le P(\sigma, e)$ .

End of the construction.

Claim 2. Consider an index e and a step  $\sigma_0$  such that

$$\forall \sigma \ge \sigma_0 \,\forall z ((P(z) \le P(e) \to (P(\sigma, z) = P(z) \land z < \sigma_0)) \land (P(z) \ge P(e) \to P(\sigma, z) \ge P(e))).$$

Define  $K_e := \{x \mid P(x) < P(e)\}$ . Then the following holds: If  $\sigma > \sigma_0$  is a step such that  $R^{K_e} \cap \sigma = R_{\sigma}^{K_e} \cap \sigma$  and no element of  $R^{K_e} \cap \sigma$  is put into A at step  $\sigma$  then no  $x < r(e, \sigma)$  is put into A at any stage  $\sigma' \ge \sigma$  and we have  $\forall \sigma' \ge \sigma(r(e, \sigma') \ge r(e, \sigma))$ .

**Proof.** Induction on P(e). Assume for a contradiction that there exists a minimal  $\sigma' \ge \sigma$  such that some  $x < r(e, \sigma)$  is put into A at step  $\sigma'$ .

Then we have  $\sigma' > \sigma$  because for  $\sigma' = \sigma$  the ordinal  $x < r(e, \sigma)$  would be an element of  $R^{K_e} \cap \sigma$ . The computations in A which contribute to the definition of  $r(e, \sigma)$  are not destroyed before  $\sigma'$  because of the minimality of  $\sigma'$ . Therefore we have  $r(e, \sigma') \ge r(e, \sigma)$  so that the element  $x = \langle x', e' \rangle < r(e, \sigma)$  which is put into A at step  $\sigma'$  is an element of  $R^{K_e} \cap \sigma$  and we have P(e') < P(e). Further since x is not put into A at step  $\sigma$  there exists some  $\hat{e}$  with  $P(\hat{e}) \le P(e') < P(e)$  such that  $x < r(\hat{e}, \sigma)$ . But this situation is impossible according to the induction hypothesis for  $P(\hat{e})$ . Contradiction.

It remains to show that  $r(e, \sigma') \ge r(e, \sigma)$  for  $\sigma' \ge \sigma$  but this follows immediately from the preceding because the computations in A which contribute to the definition of  $r(e, \sigma)$  are never destroyed.

Claim 3. For every  $e \in \alpha$  we have  $A^{K_e} = R^{K_e}$  where  $K_e := \{i \mid P(i) < P(e)\}$  and  $M_1 = M_2 : \Leftrightarrow M_1 - M_2 \in L_{\alpha} \land M_2 - M_1 \in L_{\alpha}$ .

**Proof.** Induction on P(e). We define for the proof

$$R(e, \sigma) := \sup \{ r(i, \sigma) \mid i \leq \sigma \land P(\sigma, i) \leq P(\sigma, e) \}.$$

Further we fix a step  $\sigma_0$  such that

$$\forall \sigma \ge \sigma_0 \ \forall z ((P(z) \le P(e') \to (P(\sigma, z) = P(z) \land z < \sigma_0)) \land (P(z) > P(e') \to P(\sigma, z) > P(e')).$$

Case 1. P(e') = P(e) + 1 for some e'.

We show that there exists a step  $\sigma_1$  and a constant  $r_1$  such that

$$\forall \sigma \geq \sigma_1(R(e,\sigma) \geq r_1 \land \exists \tau \geq \sigma(R(e,\tau) = r_1)).$$

By the induction hypothesis we have  $A^{K_e} = {}^{*}R^{K_e}$ . Therefore  $A^{K_e}$  is  $\alpha$ -recursive (use Claim 1)) and regular. Thus the set

$$M_e := \{ \sigma \ge \sigma_0 \mid A^{K_e} \cap \sigma = A^{K_e}_{\sigma^e} \cap \sigma \}$$

of fixpoints is unbounded in  $\alpha$  and  $\alpha$ -recursive.

A computation in A which contributes to the definition of  $R(e, \sigma)$  for some  $\sigma \ge \sigma_0$  can only be destroyed later by some element of  $A^{K_e}$  which is enumerated into A. This implies that for  $\sigma \in M_e$  the computations in A which contribute to the definition of  $R(e, \sigma)$  will never be destroyed.

Define for  $i \in K_{e'}$   $l(i) := \sup \{l(i, \sigma) \mid \sigma \in M_e\}$  and

$$H_{i} := \begin{cases} \{x+1 \mid x < l(i)\} & \text{if } \forall \sigma \ge \sigma_{0}(\sigma \in M_{e} \rightarrow q(i, \sigma) = 0), \\ \{0\} \cup \{x+1 \mid x < l(i)\} & \text{otherwise.} \end{cases}$$

Then  $K := \bigcup \{H_i \times \{i\} \mid i \in K_{e'}\}$  is  $\alpha$ -finite because K is  $\alpha$ -r.e.,  $\alpha$ -card  $(K_{e'}) < \sigma^2$  of  $\alpha$  and every  $H_i$  is  $\alpha$ -finite.

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We define an  $\alpha$ -recursive function  $f: K \rightarrow \alpha$  such that

$$f(x) = \begin{cases} \mu \sigma \in M_e(q(i, \sigma) \neq 0) & \text{if } x = \langle 0, i \rangle, \\ \mu \sigma \in M_e(l(i, \sigma) > x') & \text{if } x = \langle x' + 1, i \rangle. \end{cases}$$

f is in fact  $\alpha$ -finite because dom f = K is  $\alpha$ -finite.

Choose  $\sigma_1 \in M_e$  such that  $\operatorname{Rg} f \subseteq \sigma_1$ . Define  $r_1 := R(e, \sigma_1)$ . Then  $\sigma_1, r_1$  have the desired properties.

Choose  $r_2 \ge r_1$ ,  $\sigma_1$  such that  $R^{(e)} \cap r_1 = R^{(e)}_{r_2} \cap r_1$  and  $r_2 \in M_e$ . Then we have  $A^{(e)} \cap r_1 = A^{(e)}_{r_2+1} \cap r_1 \in L_{\alpha}$  and  $A^{(e)} - r_1 = R^{(e)} - r_1$  by the properties of  $\sigma_1, r_1$ .

Case 2. P(e') is a limit ordinal.

In this case we have  $\omega < \sigma 2$  cf  $\alpha$ . Therefore the set

$$M_{e'} := \{ \sigma > \sigma_0 \mid R^{K_{e'}} \cap \sigma = R^{K_{e'}}_{\sigma} \cap \sigma \wedge A^{K_{e'}} \cap \sigma = A^{K_{e'}}_{\sigma} \cap \sigma \}$$

is unbounded in  $\alpha$ .

Further consider the set

$$N_{e'} := \{ \sigma > \sigma_0 \mid R^{K_{e'}} \cap \sigma = R^{K_{e'}}_{\sigma} \cap \sigma \land \text{(no element } x \in R^{K_{e'}} \cap \sigma \text{ is put into } A \text{ at step } \sigma \text{)} \}.$$

 $N_{e'}$  is  $\alpha$ -recursive by Claim 1 and by using Claim 2 we get that  $M_{e'} = N_{e'}$ . Therefore  $M_{e'}$  is  $\alpha$ -recursive which implies that  $A^{K_{e'}}$  is  $\alpha$ -recursive as well.

Thus the set  $H := R^{K_{e'}} - A^{K_{e'}}$  is  $\alpha$ -r.e. Since by the induction hypothesis every  $H^{(i)}$  is  $\alpha$ -finite for  $i \in K_{e'}$  and since  $\alpha$ -card  $(K_{e'}) < \sigma^2$  cf  $\alpha$  we get that H is in fact  $\alpha$ -finite.

Claim 4. A is hyperregular.

**Proof.** Assume for a contradiction that  $\rho = \operatorname{rcf} A < \alpha$ . We have then  $\rho \le \alpha^*$  since  $\alpha^*$  is the greatest  $\alpha$ -cardinal if  $\alpha^* < \alpha$ .

There is an index  $i \in \alpha$  such that  $\{i\}^A$  is a cofinal function from  $\rho$  into  $\alpha$  and such that

$$\forall \gamma < \rho \; \exists \sigma' \; \forall \sigma \geq \sigma' \; \forall x \leq \gamma(\{i\}_{\sigma^{\sigma}}^{A}(x)\downarrow).$$

We consider then  $e := \langle i, \rho \rangle$  and the  $\alpha$ -finite set

$$K_e := \{ y \mid P(y) < P(e) \}.$$

Fix a step  $\sigma_0$  such that

$$\forall \sigma \ge \sigma_0 \,\forall z ((P(z) \le P(e) \rightarrow (P(\sigma, z) = P(z) \land z < \sigma_0)) \\ \land (P(z) \ge P(e) \rightarrow P(\sigma, z) \ge P(e))).$$

 $A^{K_{e}}$  is  $\alpha$ -recursive (and therefore regular) according to Claim 3. Therefore the set

$$M_e := \{ \sigma > \sigma_0 \mid A^{K_e} \cap \sigma = A^{K_e}_{\sigma} \cap \sigma \}$$

is unbounded in  $\alpha$  and  $\alpha$ -recursive.

We have then for every  $x \in \rho$  and  $y \in \alpha$ :

$$\{i\}^{A}(x) \simeq y \Leftrightarrow \exists \sigma \in M_{e}(1(e, \sigma) > x \land \{i\}^{A}_{\sigma}(x) \simeq y).$$

Thus we have found an  $\alpha$ -recursive definition of the cofinal function  $\{i\}^A$  which contradicts the admissibility of  $\alpha$ .

Claim 5.  $A' \leq_{\alpha} D$ .

**Proof.** A' is tame  $\Sigma_2 L_{\alpha}$  since A is hyperregular and  $\alpha$ -r.e. Therefore " $K \in L_{\alpha} \wedge K \subseteq A'$ " is a  $\Sigma_2 L_{\alpha}$  fact and since  $0' \leq_{\alpha} D$  we can express this fact  $\alpha$ -recursively in D.

For the other part of the reduction we observe that " $K \subseteq L_{\alpha} - A$ " is a  $\Pi_1 \langle L_{\alpha}, A \rangle$  fact. Therefore it is enough to express  $e \notin A'$   $\alpha$ -recursively in D. We have

$$e \notin A' \leftrightarrow \exists p \gamma \sigma_0(\gamma = P(e) \land p = P(\sigma_0, \cdot)^{-1} \upharpoonright (\gamma + 1)$$
  
 
$$\land \operatorname{dom} p = (\gamma + 1) \land \forall x y \sigma((\sigma \ge \sigma_0 \land \langle x, y \rangle \in p) \rightarrow P(\sigma, y) = x)$$
  
 
$$\land \exists K_e K H \in L_\alpha(K_e = \operatorname{Rg}(p \upharpoonright \gamma) \land R^{K_e} = H \cup \alpha \times K$$
  
 
$$\land \forall \sigma \ge \sigma_0((R^{K_e}_{\sigma^e} \cap \sigma = (H \cup \alpha \times K) \cap \sigma \land (\text{no element of } R^{K_e}_{\sigma^e} \cap \sigma$$
  
is put into A at step  $\sigma)) \rightarrow q(e, \sigma) = 0))).$ 

The existence of the  $\alpha$ -finite sets K, H on the right side was shown in Claim 1 by using the regularity of  $\hat{D}$ . It is easy to see that one can express " $R^{K_e} = H \cup \alpha \times K$ "  $\alpha$ -recursively in D. The  $\alpha$ -finite function p on the right side is only mentioned in order to be able to describe the properties of  $\sigma_0$ .  $\sigma_0$  plays a similar role as in the preceding claims.

The proof of the equivalence is then immediate from the following observation: Take  $\gamma := P(e)$  and let  $\sigma_0$  be a step such that

$$\forall \sigma \geq \sigma_0 \forall z \leq \gamma (P(\sigma, \cdot)^{-1}(z) \downarrow \land P(\sigma, \cdot)^{-1}(z) = P^{-1}(z)).$$

Define  $K_e := \{y \mid P(y) < P(e)\}$ . Define

$$M_e := \{ \sigma \ge \sigma_0 \mid R_{\sigma}^{K_e} \cap \sigma = R^{K_e} \cap \sigma \land (\text{no element of } R^{K_e} \cap \sigma \text{ is put} \\ \text{into } A \text{ at step } \sigma ) \}.$$

Since  $R^{K_e}$  and  $A^{K_e}$  are  $\alpha$ -recursive according to Claim 3 there exists an unbounded set of steps  $\sigma$  where

$$R_{\sigma^{e}}^{K_{e}} \cap \sigma = R^{K_{e}} \cap \sigma \wedge A_{\sigma^{e}}^{K_{e}} \cap \sigma = A^{K_{e}} \cap \sigma.$$

Therefore the set  $M_e$  is unbounded in  $\alpha$  which implies that

 $\forall \sigma \in M_e(q(e, \sigma) = 0) \rightarrow e \notin A'$ 

(use the regularity of A). Further Claim 2 implies that

 $\exists \sigma \in M_e(q(e, \sigma) > 0) \rightarrow e \in A'.$ 

Thus we have expressed  $e \notin A' \alpha$ -recursively in D. Claim 6.  $D \leq_{\alpha} A'$ .

**Proof.** We show that  $\hat{D} \leq_{\alpha} A'$ . According to the choice of  $\hat{D}$  it is enough to show that  $\hat{D} \leq_{w\alpha} A'$ .

The part  $e \in \hat{D}$  is a  $\Sigma_2 L_{\alpha}$  fact and therefore trivially  $\alpha$ -recursive in A' because of  $0' \leq_{\alpha} A'$ .

Concerning the part  $e \notin \hat{D}$  we have according to Claim 3

$$e \notin \hat{D} \leftrightarrow^{|} \exists \gamma (\neg \exists \delta > \gamma(\langle \delta, e \rangle \notin A)) \leftrightarrow \exists \gamma(\langle p, \langle \gamma, e \rangle) \notin A')$$

for some fixed parameter p.  $\Box$  –

This finishes the proof of Theorem 4.

It is tempting to define — in analogy to the definition of high and low degrees — for  $\alpha$ -degrees **a**:

**a** is intermediate :  $\Leftrightarrow \mathbf{a}' = 0^{\frac{3}{2}}$ .

 $0^{\frac{3}{2}}$  does not exist for every admissible  $\alpha$ . Therefore we consider as well the following definition which makes sense for every  $\alpha$ :

**a** is intermediate:  $\Leftrightarrow \mathbf{a}'$  is equal to the greatest  $\Delta_2 L_{\alpha}$  degree and to the greatest tame  $\Sigma_2 L_{\alpha}$  degree.

According to this second definition intermediate  $\alpha$ -r.e. degrees exist exactly for those  $\alpha$  where incomplete non-hyperregular  $\alpha$ -r.e. degrees exist. Since for these  $\alpha$  the degree  $0^{\frac{3}{2}}$  is well defined we see that both definitions characterize the same class of  $\alpha$ -r.e. degrees.

By the preceding theorems there exist hyperregular and non-hyperregular intermediate  $\alpha$ -r.e. degrees. Further results are needed in order to see whether intermediate  $\alpha$ -r.e. degrees are relevant for the fine structure of  $\alpha$ -r.e. sets and degrees. In particular it would be nice to find intrinsic properties of the intermediate  $\alpha$ -r.e. degrees (or of those  $\alpha$ -r.e. degrees **a** which satisfy  $\mathbf{a}' \leq 0^{\frac{3}{2}}$  respectively  $\mathbf{a}' \geq 0^{\frac{3}{2}}$ ) which don't mention the jump. We get some first results in this direction by using recent work of A. Leggett [3].

Leggett shows that in the case  $\sigma 1p \alpha = \omega$  and  $\alpha$ -r.e. degree a contains a maximal set iff  $U_{2^{\alpha}}^{L_{\alpha}} \leq_{w\alpha} a'$  (where  $U_{2^{\alpha}}^{L_{\alpha}}$  is an universal  $\Sigma_2 L_{\alpha}$  set). If  $\sigma 1p \alpha = \omega$  we can write  $0^{\frac{3}{2}} \leq a'$  instead of the latter condition according to [6].

Further Leggett shows in [3] for the larger class of  $\alpha$  with  $\sigma 2p \alpha = \omega$  that an  $\alpha$ -r.e. degree **a** satisfies  $U_{2^{\alpha}}^{L_{\alpha}} \leq_{w\alpha} a'$  iff **a** contains a maximal set or **a** is non-hyperregular. Thus by the preceding we get for those  $\alpha$  with  $\sigma 2p \alpha = \omega$  which are not  $\Sigma_2$  admissible that a hyperregular  $\alpha$ -r.e. degree **a** contains a maximal set iff it is intermediate. Theorem 4 shows that hyperregular intermediate  $\alpha$ -r.e. degrees exist for all these  $\alpha$  and so Martin's Theorem fails in these cases.

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Simpson conjectured in his thesis [13] that for those  $\alpha$  where maximal sets exists we can find in fact hyperregular ones. His conjecture is now proved for those  $\alpha$  which satisfy  $\sigma 2p \alpha = \omega$ .

The final theorem characterizes those  $\alpha$  which are  $\Sigma_2$  admissible in terms of their degree structure. The proof is based on non-trivial results from  $\beta$ -recursion theory.

**Theorem 5.** Assume  $\alpha$  is admissible. Then  $\alpha$  is  $\Sigma_2$  admissible if and only if every  $\Sigma_2 L_{\alpha}$  degree  $d \ge 0'$  is the jump of an incomplete  $\alpha$ -r.e. degree.

**Proof.** If  $\sigma 2 \operatorname{cf} \alpha < \sigma 2p \alpha$ , then 0" is not the jump of an incomplete  $\alpha$ -r.e. degree according to Theorem 2 in [6].

If  $\sigma 2p \alpha \leq \sigma 2$  cf  $\alpha < \alpha$  (i.e.  $\alpha$  is of type (2)) we consider the weakly inadmissible structure  $\mathfrak{B} := \langle L_{\alpha}, C \rangle$  where  $C \in 0'$  is  $\alpha$ -r.e. and regular. By Theorem 2 there exists a  $\mathfrak{B}$ -r.e. degree d which is not  $\mathfrak{B}$ -tame r.e. such that  $d <_{\mathfrak{B}} r$  (where r is the greatest  $\mathfrak{B}$ -recursive degree). Therefore there exists a  $\Sigma_2 L_{\alpha}$  degree d which is not tame  $\Sigma_2 L_{\alpha}$  such that  $0' <_{\alpha} d <_{\alpha} 0^{\frac{3}{2}}$ . If a is an incomplete  $\alpha$ -r.e. degree then we have  $0^{\frac{3}{2}} \leq a'$  if a is non-hyperregular and a' is a tame  $\Sigma_2 L_{\alpha}$  degree if a is hyperregular. Therefore we have  $d \neq a'$  for every incomplete  $\alpha$ -r.e. degree a.

It remains to show that for  $\Sigma_2$  admissible  $\alpha$  every  $\Sigma_2 L_{\alpha}$  degree  $d \ge 0'$  is the jump of an incomplete  $\alpha$ -r.e. degree.

For  $\alpha = \omega$  this is the jump theorem of Sacks [7].

For  $\alpha > \omega$  we apply a simplified version of the construction in the proof of Theorem 4 (define  $r(e, \sigma) := q(e, \sigma)$ ).

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